

On CS Undergrad Teaching and Visualization Research

Part 1 (teaching)

**A Mini-lecture on
sorting in different programming paradigms**

Part 2 (research)

**Visual Data Analysis:
real users, real data, right methods**

Michael Sedlmair

Visualization and Data Analysis group (VDA)



universität
wien

Part 1: Teaching

**A Mini-lecture on
Sorting in different programming
paradigms**

Assumptions wrt previous knowledge

- Target audience:
 - Mainly 2nd term CS students,
 - some 1st term
- Previous knowledge of
 - Programming: Java
 - Algorithms: Complexity analysis, data structures, recursion
- Learning goals
 - Repetition/introduction of sorting algorithms
 - First contact with programming paradigms, specifically imperative VS. declarative programming

Outline

- Sorting algorithms
 - Selection Sort
 - Quicksort
- Sorting in different programming paradigms
 - Imperative VS. declarative
 - Quicksort in Java (imperative)
 - Quicksort in Haskell (declarative/functional)

Sorting algorithms

Sorting



Sorting

- One of the standard algorithmic problems in Computer Science
- Goal: Put an unsorted list of elements u in ascending/descending order.
- Example:
[7,3,3,1,5] $\longrightarrow$ [1,3,3,5,7]
 └ element u_1 ,
 value_of(u_1) = 7
- *(Note: special case of sorting numbers in ascending order)*

Sorting algorithms

- Simple sorts
 - **Selection Sort**
 - Insertion Sort
 - Bubble Sort
 - ...
- More efficient sorts
 - **Quicksort**
 - Merge Sort
 - Heap Sort
 - Radixsort
 - ...

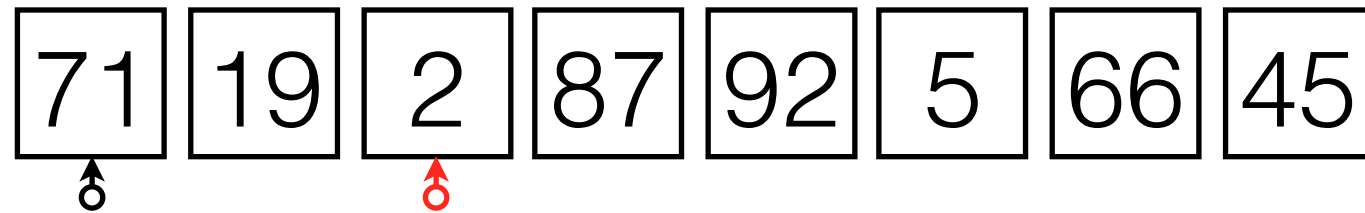
Selection Sort

- Split list into 2 parts:
 - L , sorted part of the list (left)
 - R , rest of the list (right)
- At the beginning $L = \emptyset$
- Identify the smallest element in R and swap it with the first element of R
- Done when $|R| = 1$

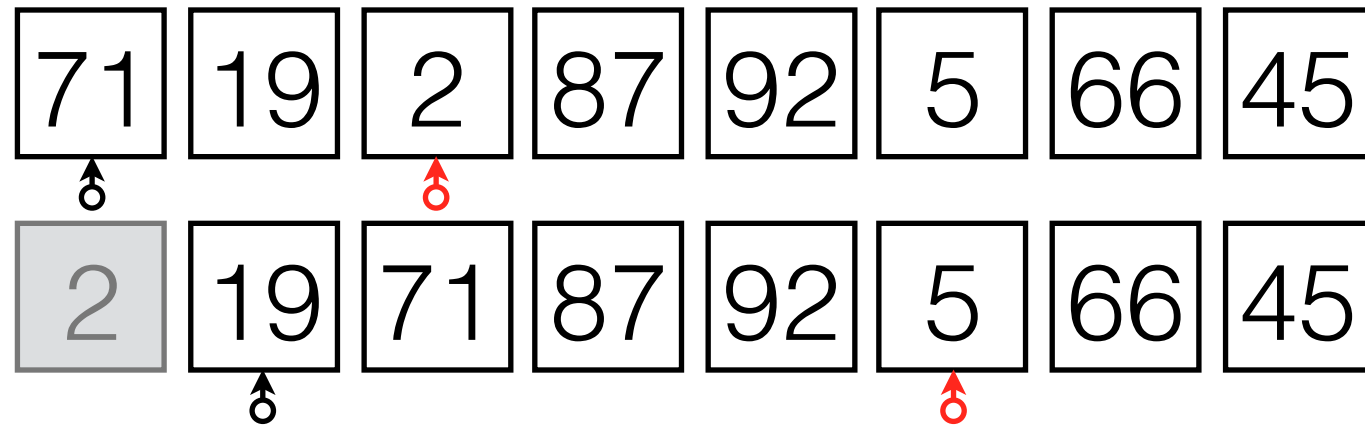
Selection Sort

71	19	2	87	92	5	66	45
----	----	---	----	----	---	----	----

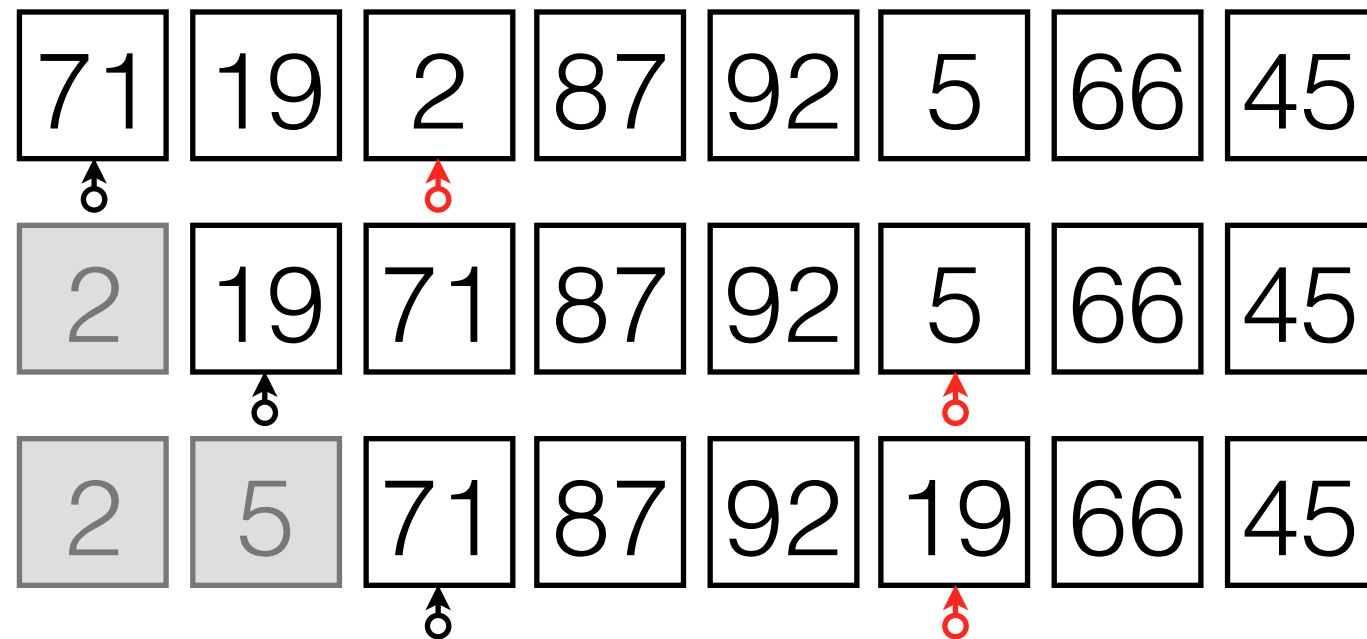
Selection Sort



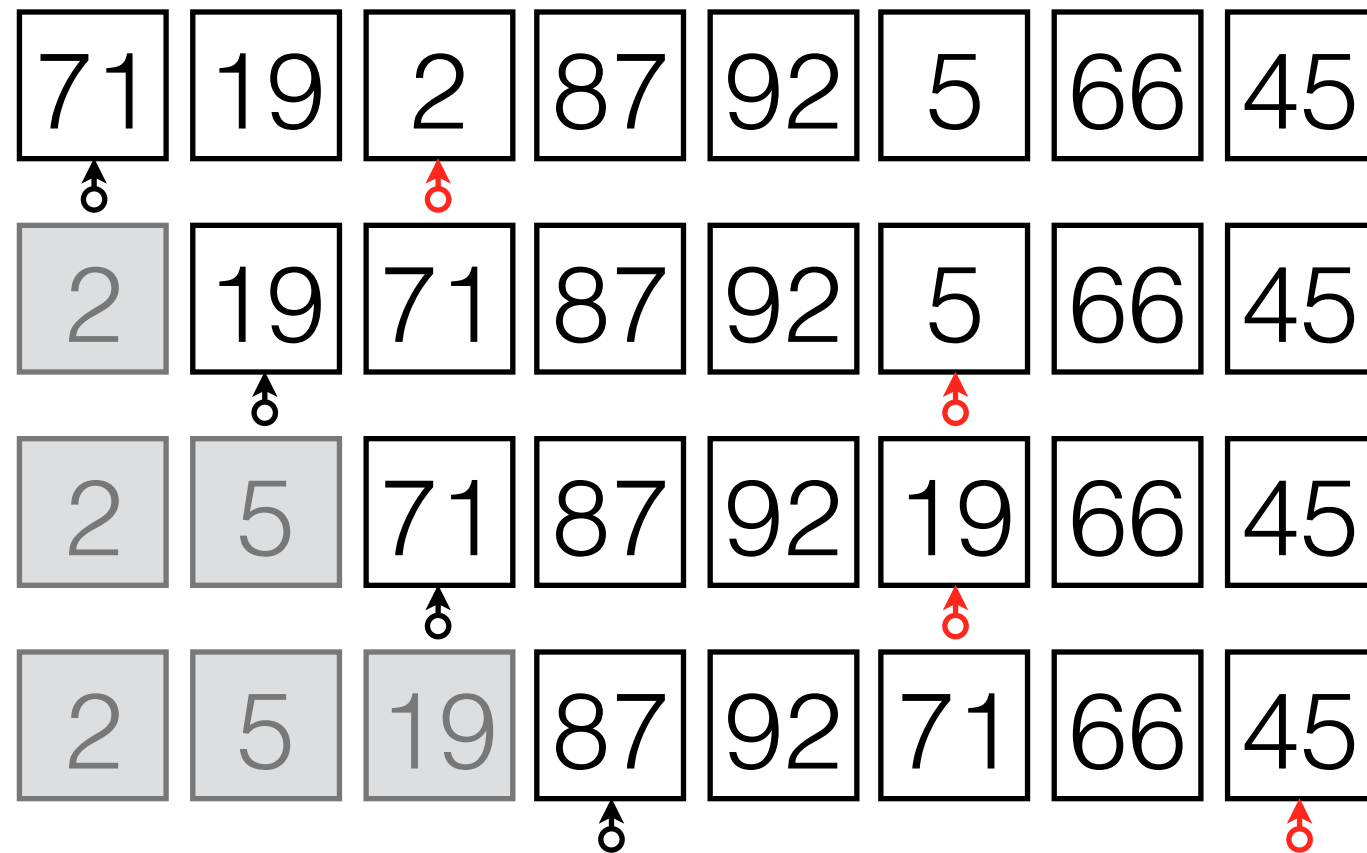
Selection Sort



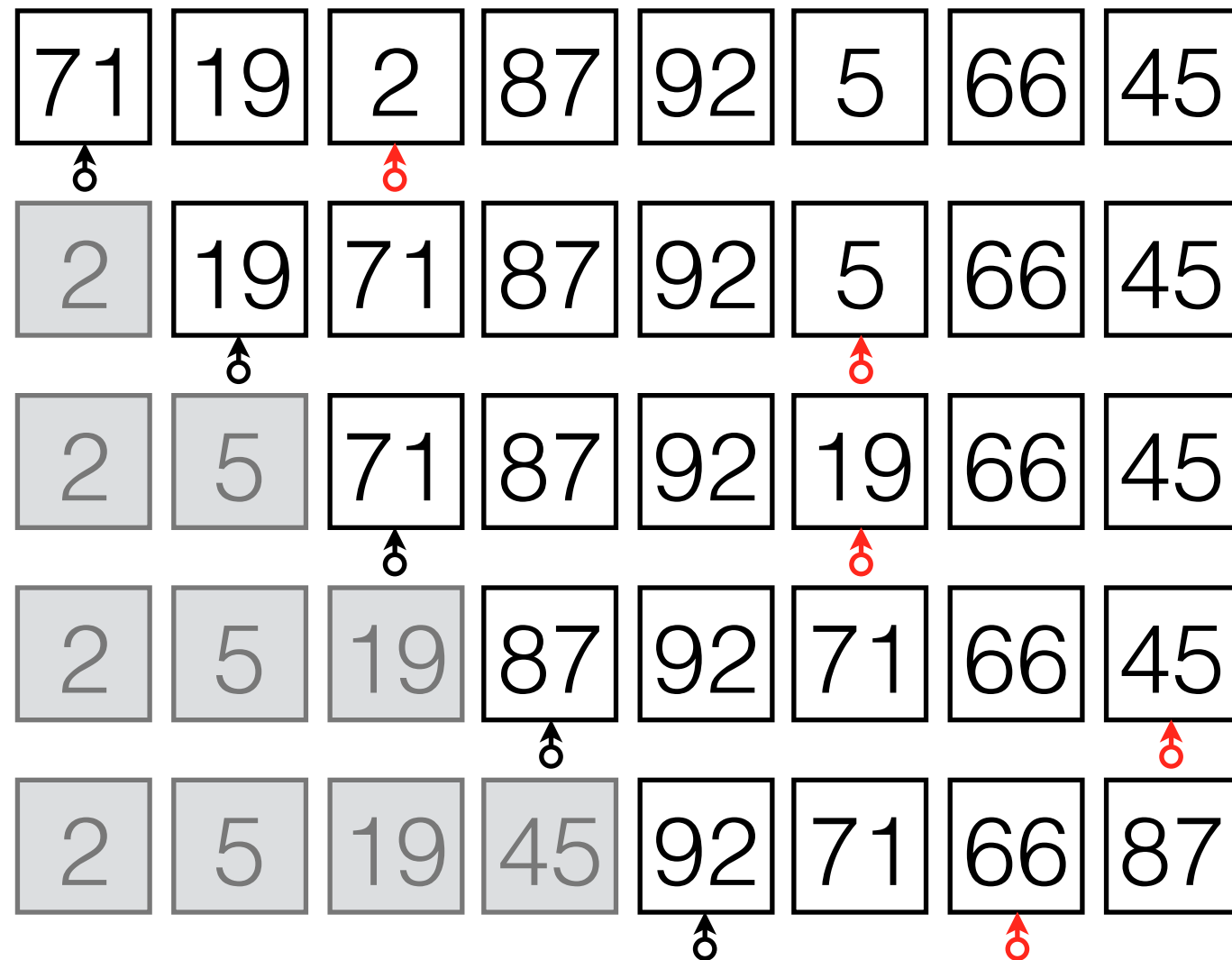
Selection Sort



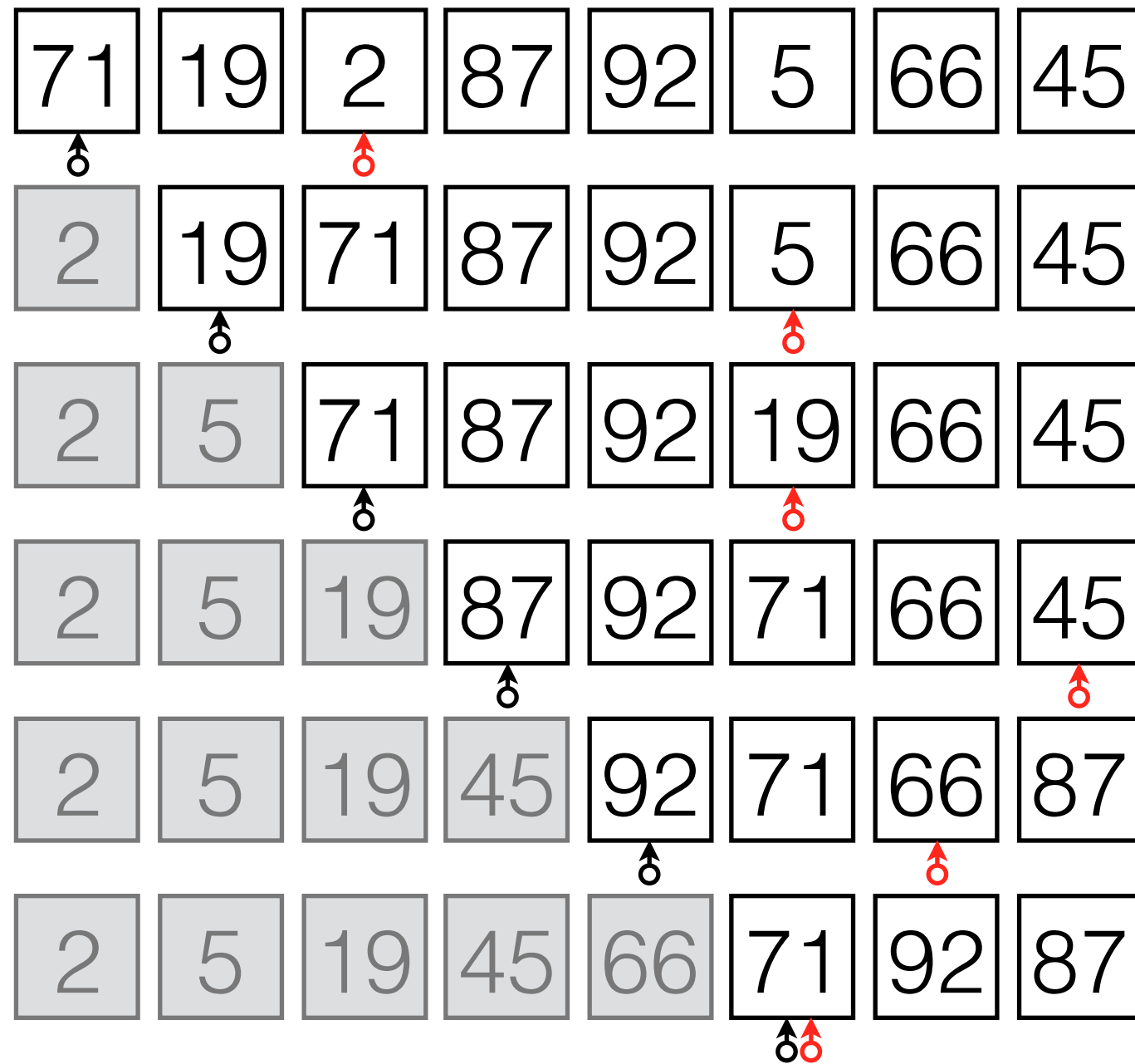
Selection Sort



Selection Sort



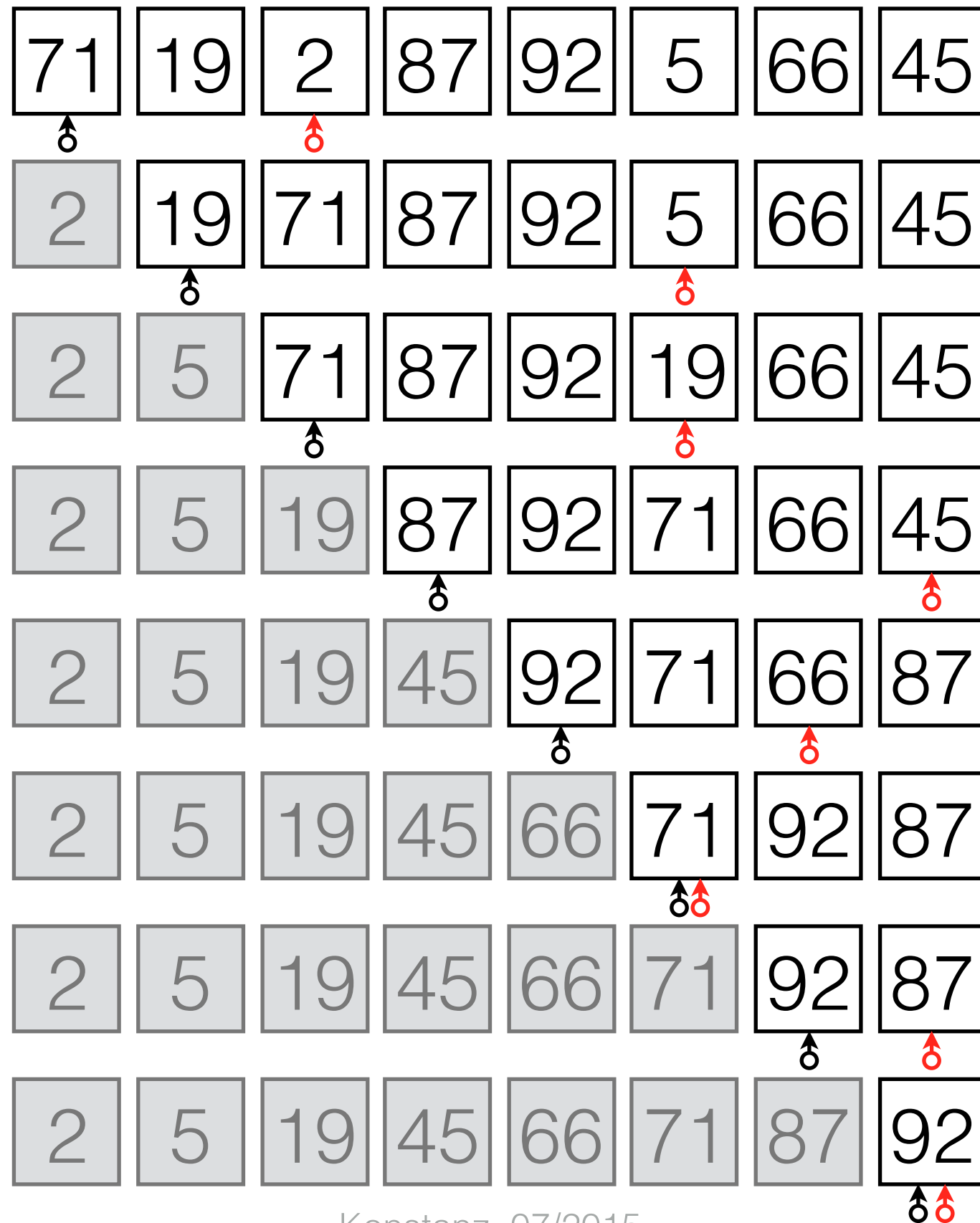
Selection Sort



Selection Sort



Selection Sort



Selection Sort: Strengths / Weaknesses

- Strengths
 - easy to understand
 - “in-place” (memory usage: n)

Selection Sort: Strengths / Weaknesses

- Strengths
 - easy to understand
 - “in-place” (memory usage: n)
- Biggest Weakness
 - Expensive for large lists: $O(n^2)$

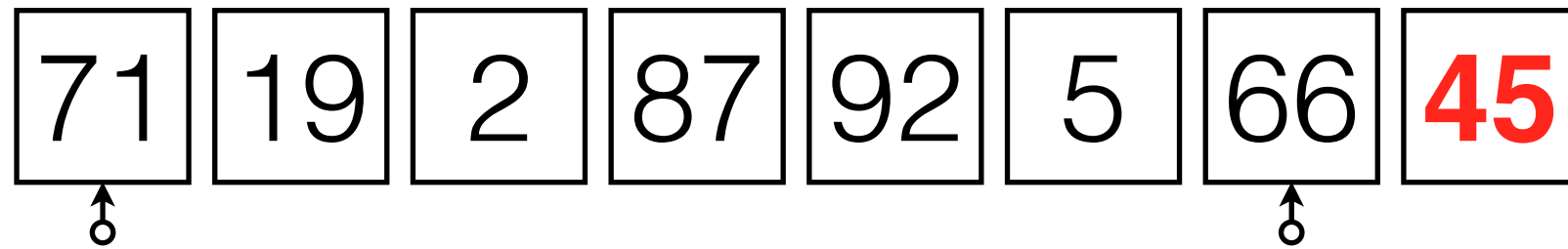
Quicksort

- Divide-and-conquer algorithms
- Determine an arbitrary *pivot element* u_p (e.g., last element of a list)
- Swap elements in the list, so that
 - all elements left of u_p are smaller than u_p
 - all elements right of u_p are greater than u_p
- The *pivot* u_p is then already at its final position
- Recursively sort left and right parts of the list

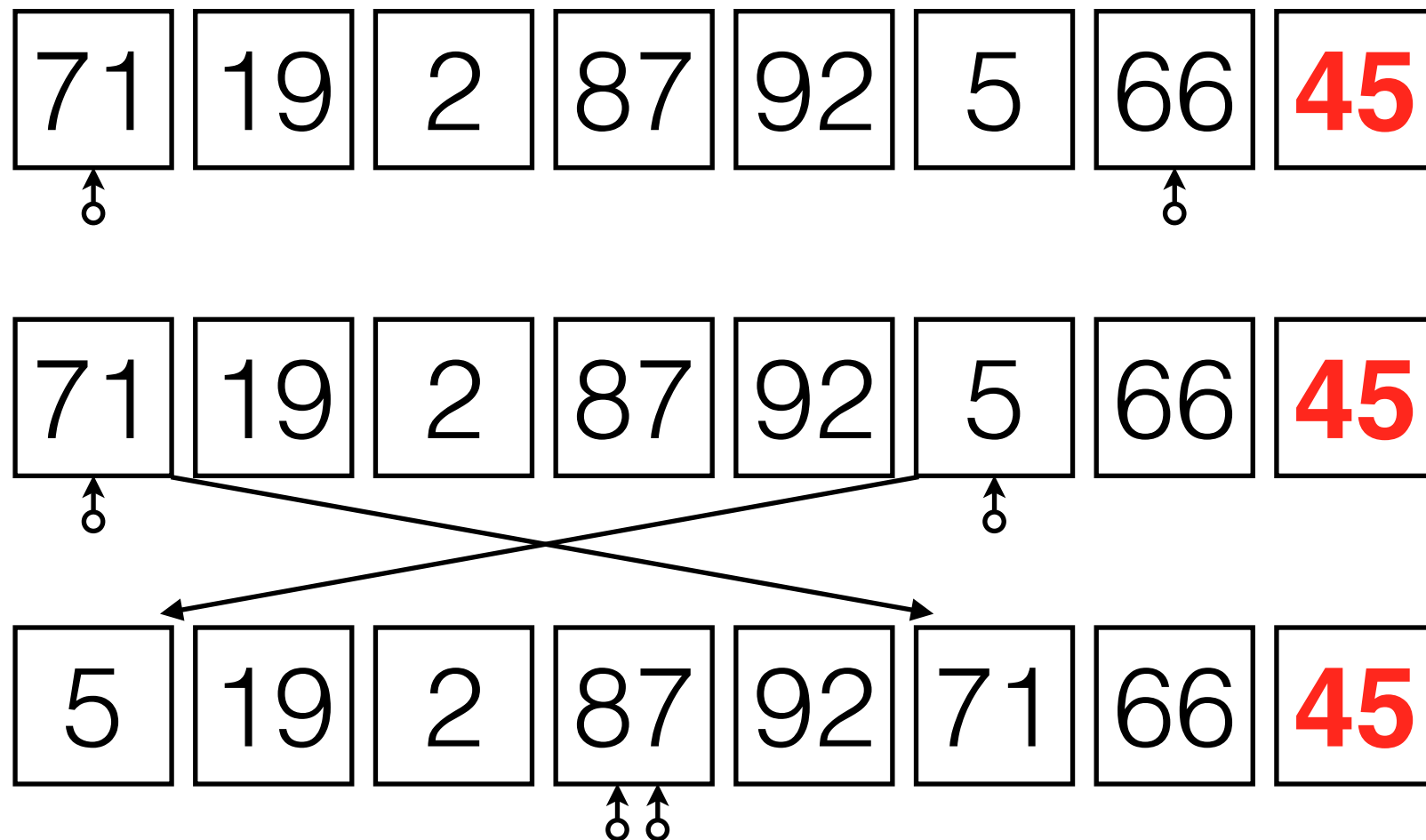
Quicksort

- Moving the pivot:
 - Go from left to right until you find an element u_i with $\text{value_of}(u_i) > \text{value_of}(u_p)$
 - Go from right to left until you find an element u_j with $\text{value_of}(u_j) < \text{value_of}(u_p)$
 - Swap these elements
 - Repeat until i and j meet
 - Swap *pivot element* with the element at position i if $\text{value_of}(u_p) < \text{value_of}(u_i)$

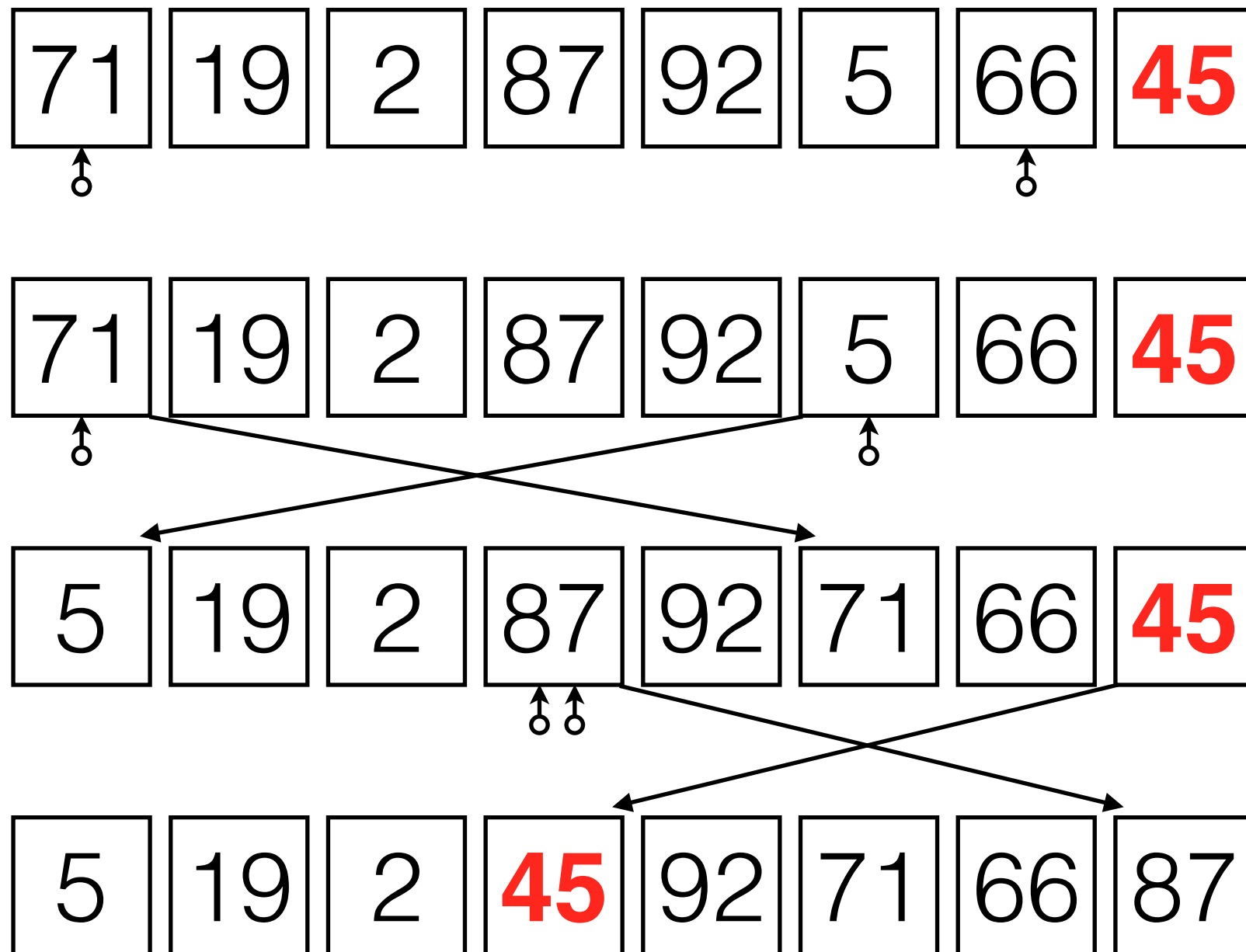
Quicksort



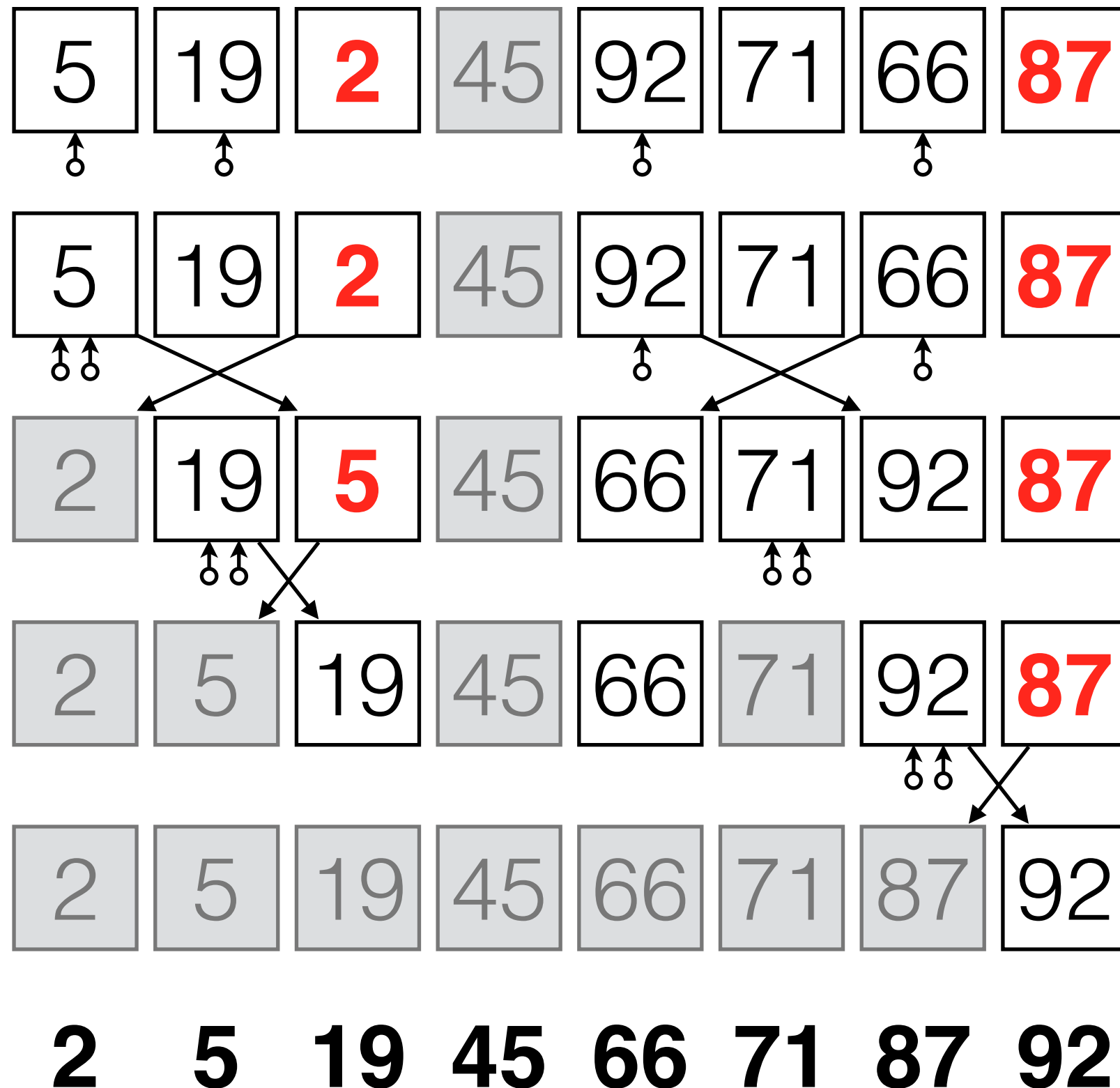
Quicksort



Quicksort



Quicksort



Quicksort: Strength / Weaknesses

- Strength:
 - Average complexity: $O(n \cdot \log_2(n))$
 - “in-place” memory

Quicksort: Strength / Weaknesses

- Strength:
 - Average complexity: $O(n \cdot \log_2(n))$
 - “in-place” memory
- Weaknesses:
 - Worst case complexity: $O(n^2)$
 - not efficient for small lists

Sorting in different programming paradigms

(using the example Quicksort)

Programming paradigms

- Fundamental style or “way” of programming
- Some examples
 - Imperative — Give an explicit sequence of commands
 - Declarative — Programs specify the result, not how to get it.
 - Object-Oriented — Computation is effected by sending messages to objects; objects have state and behavior.
 - ...
- Not mutually exclusive concepts! Specifically, languages can follow more than one paradigm

Imperative VS. declarative

Imperative VS. declarative

- Step-by-step recipe
- Describe the how: Control flow in imperative programming is explicit; commands show how the computation takes place.
- Examples: Java, C++, C, ...



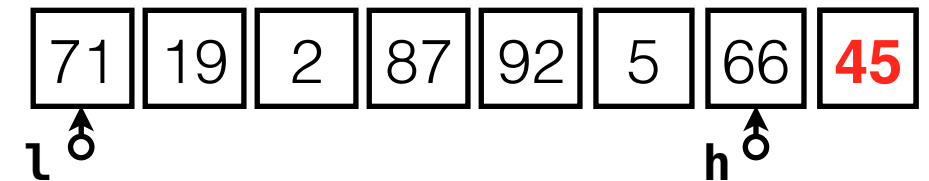
Imperativ: Quicksort in Java

```
public static void qsort (int[] val, int low, int high) {  
  
    if (high - low > 0){  
        //get pivot element from the end of the list  
        int pivot = high;  
  
        //initialize pointers  
        int l = low, h = high - 1;  
  
        while (l < h) {  
            //set pointers  
            while ((l < h) && (val[l] <= val[pivot])) {  
                l++;  
            }  
            while ((l < h) && (val[h] >= val[pivot])) {  
                h--;  
            }  
  
            //swap values  
            if (l != h) {  
                swap(val, l, h);  
            } else if (val[l] > val[pivot]) {  
                swap(val, l, pivot);  
            }  
        }  
        //recursion  
        quicksort(val, low, l - 1);  
        quicksort(val, l + 1, high);  
    }  
}
```

71	19	2	87	92	5	66	45
----	----	---	----	----	---	----	----

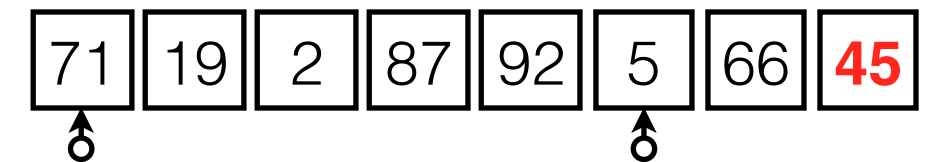
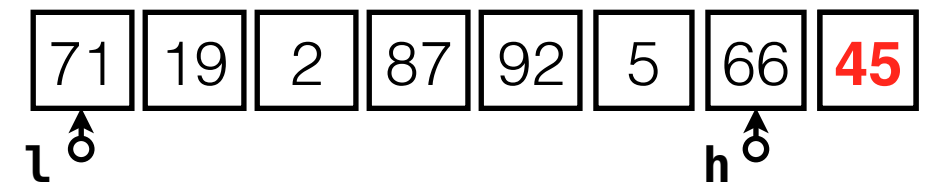
Imperativ: Quicksort in Java

```
public static void qsort (int[] val, int low, int high) {  
  
    if (high - low > 0){  
        //get pivot element from the end of the list  
        int pivot = high;  
  
        //initialize pointers  
        int l = low, h = high - 1;  
  
        while (l < h) {  
            //set pointers  
            while ((l < h) && (val[l] <= val[pivot])) {  
                l++;  
            }  
            while ((l < h) && (val[h] >= val[pivot])) {  
                h--;  
            }  
  
            //swap values  
            if (l != h) {  
                swap(val, l, h);  
            } else if (val[l] > val[pivot]) {  
                swap(val, l, pivot);  
            }  
        }  
        //recursion  
        quicksort(val, low, l - 1);  
        quicksort(val, l + 1, high);  
    }  
}
```



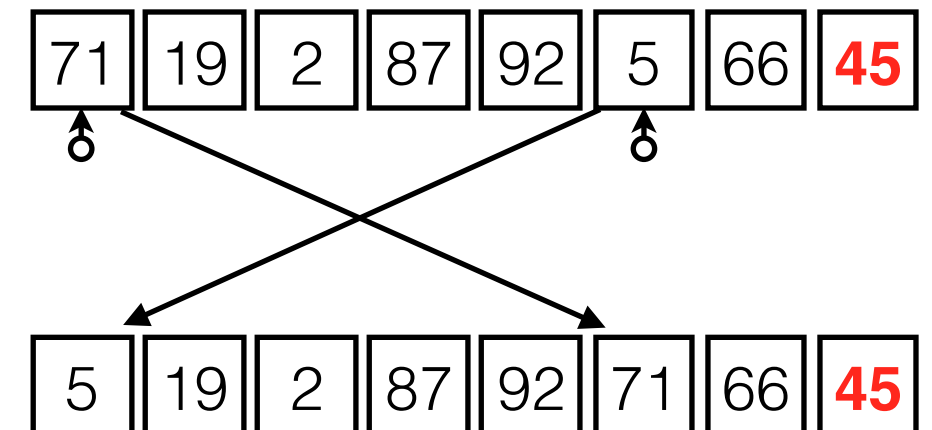
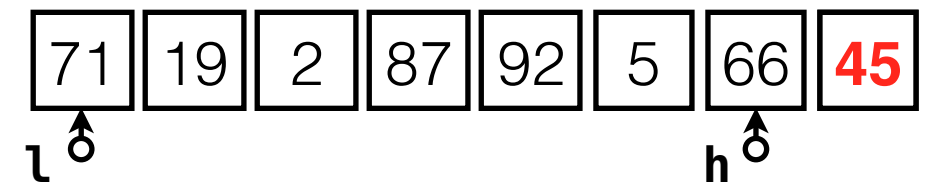
Imperativ: Quicksort in Java

```
public static void qsort (int[] val, int low, int high) {  
    if (high - low > 0){  
        //get pivot element from the end of the list  
        int pivot = high;  
  
        //initialize pointers  
        int l = low, h = high - 1;  
  
        while (l < h) {  
            //set pointers  
            while ((l < h) && (val[l] <= val[pivot])) {  
                l++;  
            }  
            while ((l < h) && (val[h] >= val[pivot])) {  
                h--;  
            }  
  
            //swap values  
            if (l != h) {  
                swap(val, l, h);  
            } else if (val[l] > val[pivot]) {  
                swap(val, l, pivot);  
            }  
        }  
        //recursion  
        quicksort(val, low, l - 1);  
        quicksort(val, l + 1, high);  
    }  
}
```



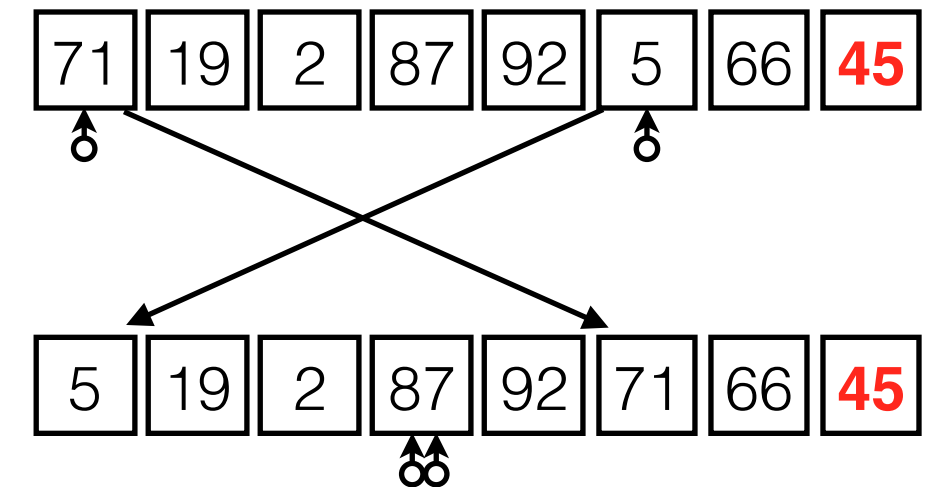
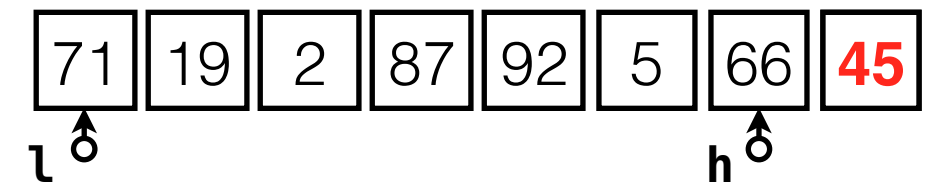
Imperativ: Quicksort in Java

```
public static void qsort (int[] val, int low, int high) {  
    if (high - low > 0){  
        //get pivot element from the end of the list  
        int pivot = high;  
  
        //initialize pointers  
        int l = low, h = high - 1;  
  
        while (l < h) {  
            //set pointers  
            while ((l < h) && (val[l] <= val[pivot])) {  
                l++;  
            }  
            while ((l < h) && (val[h] >= val[pivot])) {  
                h--;  
            }  
  
            //swap values  
            if (l != h) {  
                swap(val, l, h);  
            } else if (val[l] > val[pivot]) {  
                swap(val, l, pivot);  
            }  
        }  
        //recursion  
        quicksort(val, low, l - 1);  
        quicksort(val, l + 1, high);  
    }  
}
```



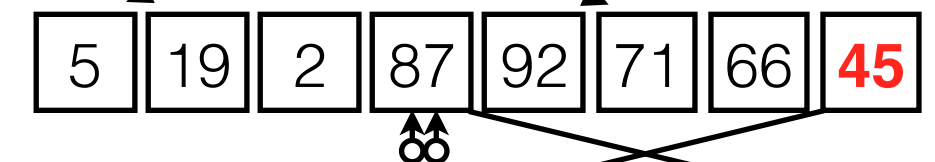
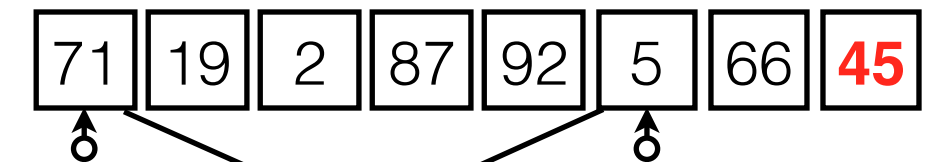
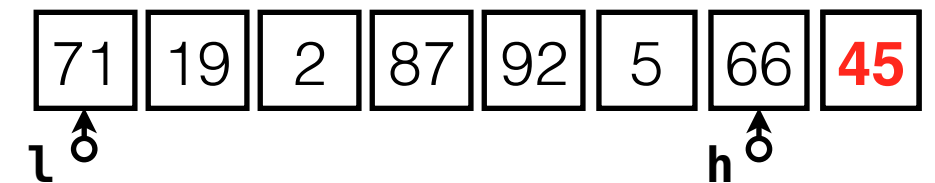
Imperativ: Quicksort in Java

```
public static void qsort (int[] val, int low, int high) {  
    if (high - low > 0){  
        //get pivot element from the end of the list  
        int pivot = high;  
  
        //initialize pointers  
        int l = low, h = high - 1;  
  
        while (l < h) {  
            //set pointers  
            while ((l < h) && (val[l] <= val[pivot])) {  
                l++;  
            }  
            while ((l < h) && (val[h] >= val[pivot])) {  
                h--;  
            }  
  
            //swap values  
            if (l != h) {  
                swap(val, l, h);  
            } else if (val[l] > val[pivot]) {  
                swap(val, l, pivot);  
            }  
        }  
        //recursion  
        quicksort(val, low, l - 1);  
        quicksort(val, l + 1, high);  
    }  
}
```



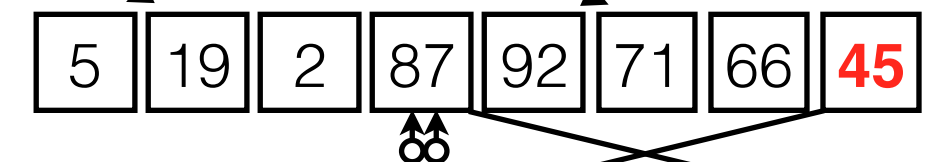
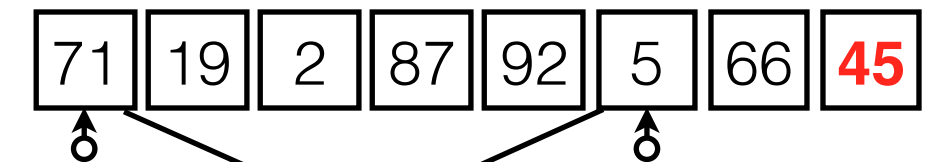
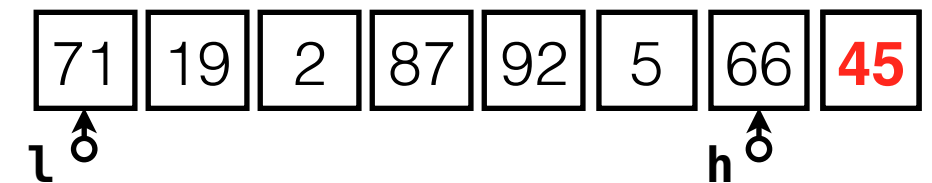
Imperativ: Quicksort in Java

```
public static void qsort (int[] val, int low, int high) {  
    if (high - low > 0){  
        //get pivot element from the end of the list  
        int pivot = high;  
  
        //initialize pointers  
        int l = low, h = high - 1;  
  
        while (l < h) {  
            //set pointers  
            while ((l < h) && (val[l] <= val[pivot])) {  
                l++;  
            }  
            while ((l < h) && (val[h] >= val[pivot])) {  
                h--;  
            }  
  
            //swap values  
            if (l != h) {  
                swap(val, l, h);  
            } else if (val[l] > val[pivot]) {  
                swap(val, l, pivot);  
            }  
        }  
        //recursion  
        quicksort(val, low, l - 1);  
        quicksort(val, l + 1, high);  
    }  
}
```



Imperativ: Quicksort in Java

```
public static void qsort (int[] val, int low, int high) {  
    if (high - low > 0){  
        //get pivot element from the end of the list  
        int pivot = high;  
  
        //initialize pointers  
        int l = low, h = high - 1;  
  
        while (l < h) {  
            //set pointers  
            while ((l < h) && (val[l] <= val[pivot])) {  
                l++;  
            }  
            while ((l < h) && (val[h] >= val[pivot])) {  
                h--;  
            }  
  
            //swap values  
            if (l != h) {  
                swap(val, l, h);  
            } else if (val[l] > val[pivot]) {  
                swap(val, l, pivot);  
            }  
        }  
  
        //recursion  
        quicksort(val, low, l - 1);  
        quicksort(val, l + 1, high);  
    }  
}
```



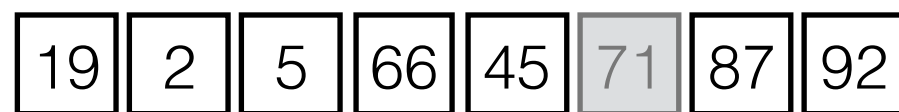
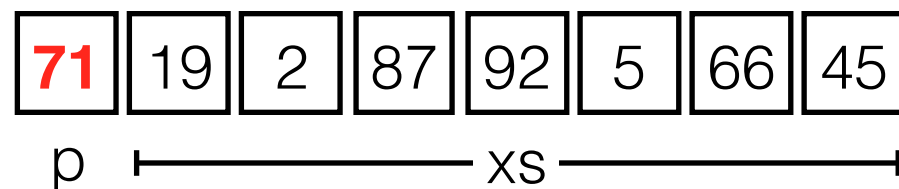
Imperative VS. declarative

- What not the how
- *“if you can describe the problem, you have solved the problem” [Brian Tarbox]*
- Higher-level abstraction
- Main benefit:
 - Code is much shorter
 - Less error-prone
- Example: Haskell (functional)



Declarative/functional: Quicksort in Haskell

```
qsort [] = []  
qsort (p:xs) = qsort lesser ++ [p] ++ qsort greater  
  where  
    lesser = [ y | y <- xs, y < p ]  
    greater = [ y | y <- xs, y >= p ]
```



- `[]` is the empty list
- `(p:xs)`
- `[p]` is a list with one element
- `++` concatenation of two list
- `lesser /greater` filters the list

Summary Mini-Lecture

- Sorting
 - One of the canonical algorithmic problems in Computer Science
 - Many algorithms exist
 - Examples: Selection Sort (simple), Quicksort (more efficient)
- Programming paradigms
 - Imperative VS. declarative
 - Quicksort in Java VS. Quicksort in Haskell

Further references

- Quicksort in 42 different programming languages:
https://en.wikibooks.org/wiki/Algorithm_Implementation/Sorting/Quicksort
- Books
 - Knuth, Donald E. (1998), Sorting and Searching, The Art of Computer Programming 3 (2nd ed.), Addison-Wesley.
 - Tate, Bruce A. (2010), Seven Languages in Seven Weeks: A Pragmatic Guide to Learning Programming Languages (1st ed.), Pragmatic Bookshelf.
- Other lectures & **special thanks (!)**
 - Konzepte der Informatik, 2013, Prof. Deussen
<https://streaming.uni-konstanz.de/lectures/wintersemester-2013/inf-11700-20132/>
 - Konzepte der Programmierung, 2014, Prof. Scholl
<https://streaming.uni-konstanz.de/lectures/sommersemester-2014/inf-12070-20141/>
 - ACM curriculum
<https://www.acm.org/education/CS2013-final-report.pdf>

Part 2: Research

**Visual Data Analysis:
real users, real data, right methods**



data

has become ubiquitous

visualization & data analysis

My Focus

- **real users/real data**

- problem-driven: design studies
- technique-driven: high-dimensional data analysis



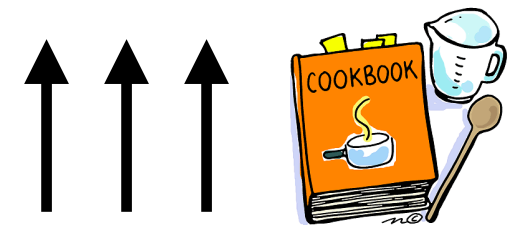
→
problem-driven



←
technique-driven

- **right methods**

- innovate and refine research methods for visual data analysis

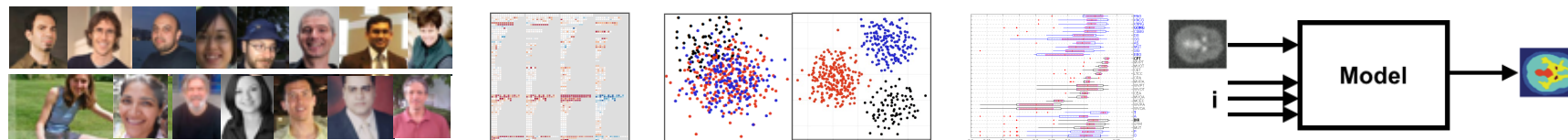


real users/real data:

Applied visualization projects (design studies)

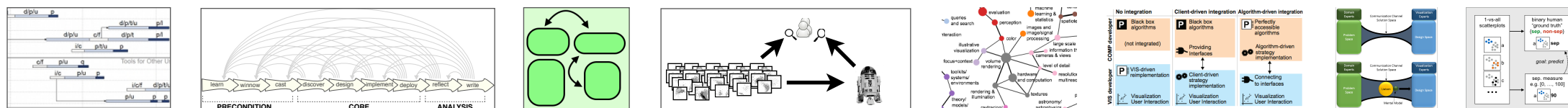


High-dimensional data analysis techniques



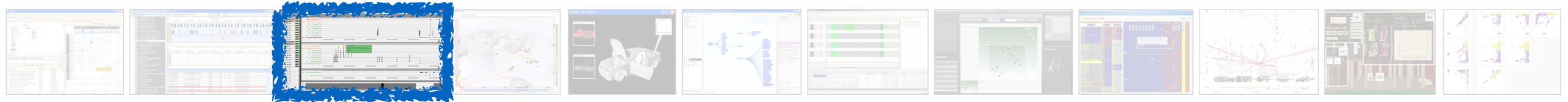
right methods:

Novel and refined research methods



real users/real data

Applied visualization projects (design studies)

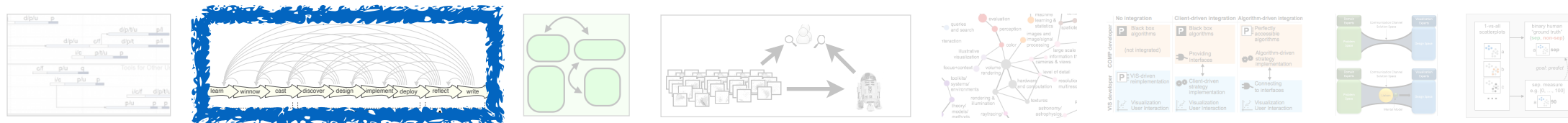


High-dimensional data analysis techniques



right methods:

Novel and refined research methods



real users/real data:

Cardiogram: Visual Analytics for Automotive Engineers

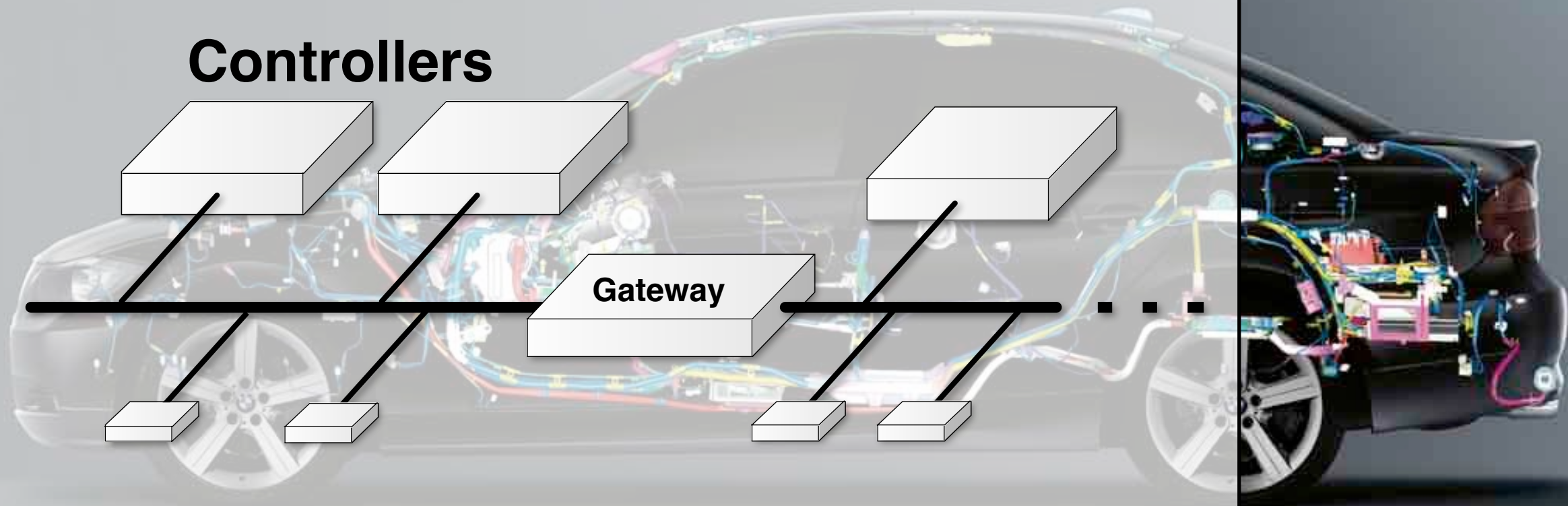
M. Sedlmair, P. Isenberg, D. Baur, M. Mauerer, C. Pigorsch, A. Butz
(*CHI 2011*).



more and more electronics...



**... enabled by in-car
communication networks**



~70 Controllers / Car

Problem Characterization



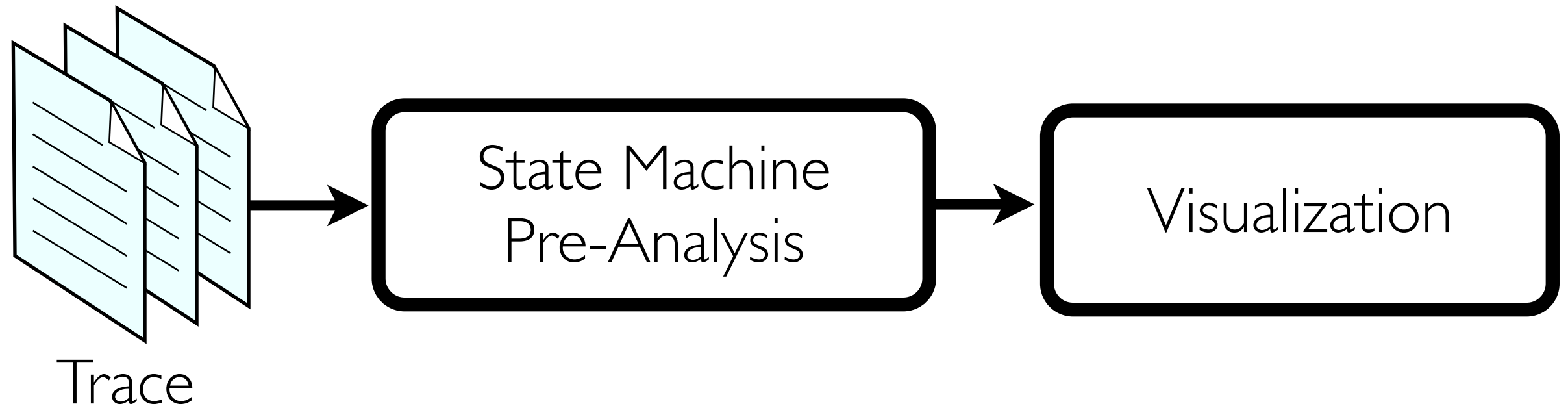
75836...	F...	55		Rx	16	5a 77 f8 27 00 20 00 20 0f 00 0...
75836...	F...	56		Rx	16	ee 7e f8 27 00 20 00 20 0f 00 0...
75836...	F...	57		Rx	16	ed 77 10 28 00 20 00 20 0f 00 0...
75836...	F...	58		Rx	16	5e f8 1f 00 02 20 ff 21 22 22 07 f...
75851...	F...	5a		Rx	16	33 08 77 22 75 d2 6f e2 70 f2 03 0...
75851...	1	1a1	V_VEH	Rx	5	c1 f7 00 00 8a
~ V_VEH_COG 0 km/h [0]						
~ ST_ECU_V_VEH Signal ungültig [f]						
~ QU_V_VEH_COG Signalwert ist gültig, Zustand/S [a]						
~ DVCO_VEH Fahrzeug steht [0]						
~ CRC_V_VEH 193 [c1]						
~ ALIV_V_VEH 7 [7]						
75851...	F...	5c		Rx	8	00 00 00 00 ff ff 00 10
75851...	1	1c4		Rx	6	00 00 00 00 ff ff
75851...	F...	12f		Rx	72	00 00 00 00 00 00 ff 59 87 21 4c 0...
75851...	1	1c5		Rx	6	02 00 04 00 ff ff
75836...	F...	1		Rx	16	10 7d 18 28 00 20 00 20 0f 00 0...
75836...	F...	2		Rx	16	ca 76 f9 27 00 20 00 20 0f 00 0...
75836...	F...	3		Rx	16	6a 7d fa 27 00 20 00 20 0f 00 0...
75836...	F...	4		Rx	16	a3 76 0e 28 00 20 00 20 0f 00 0...
75836...	F...	5		Rx	16	2c f7 1f 00 02 20 ff 21 22 22 07 f...
75851...	F...	7		Rx	16	6d f9 76 12 75 d2 6f f2 70 f2 01 1...
75851...	F...	12		Rx	0	
75851...	2	301	AVL_STEA_DV	Rx	7	51 15 f8 7f ff 7f 11
75851...	3	301	AVL_STEA_DV	Rx	7	51 15 f8 7f ff 7f 11
75851...	4	137		Rx	2	fd 00
75851...	3	d9	ANG_ACPD	Rx	8	9b 99 00 c0 00 e0 7f f0
75848...	1	299		Rx	5	9f ff ff ff ff
75851...	F...	21		Rx	0	
75851...	4	a5	TORQ_CRSH_1	Rx	8	45 f5 48 f7 7f 00 00 fc
75851...	3	a5	TORQ_CRSH_1	Rx	8	fe f9 48 f7 7f 00 00 fc
75836...	F...	23		Rx	16	7c 10 05 05 ea f3 53 20 74 10 20 f...
75851...	F...	26		Rx	0	

- **data:** recorded traces (15k messages/sec)
- **task:** finding errors
- **process:** test drives
- **current practices:** mainly textual lists

Design Requirements

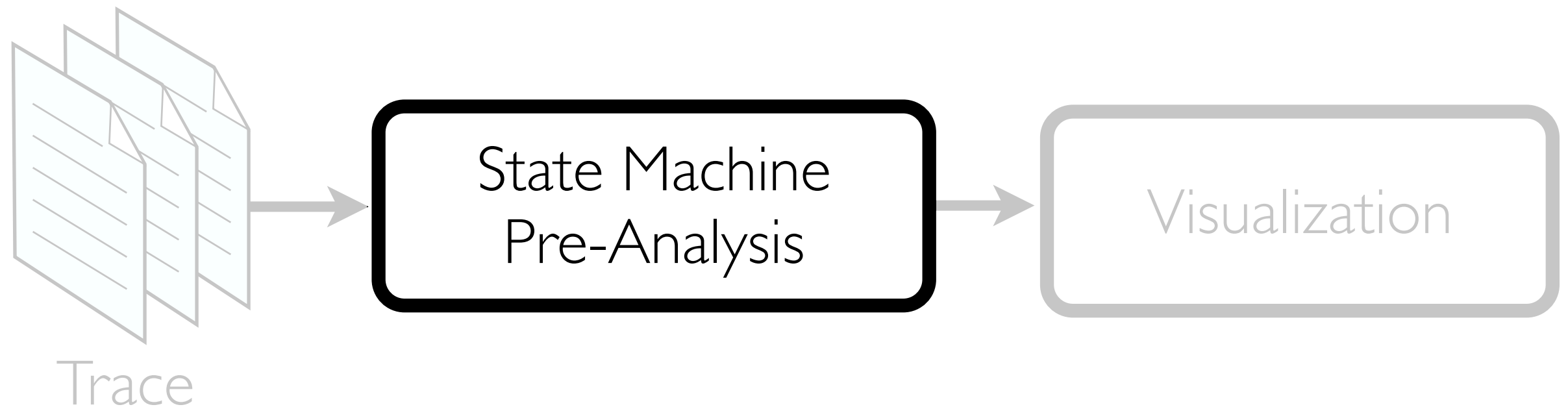
- **Main Challenges**
 - handling masses of test traces (large and many)
 - understanding correlation between trace and car behavior
 - ...

Cardiogram overview

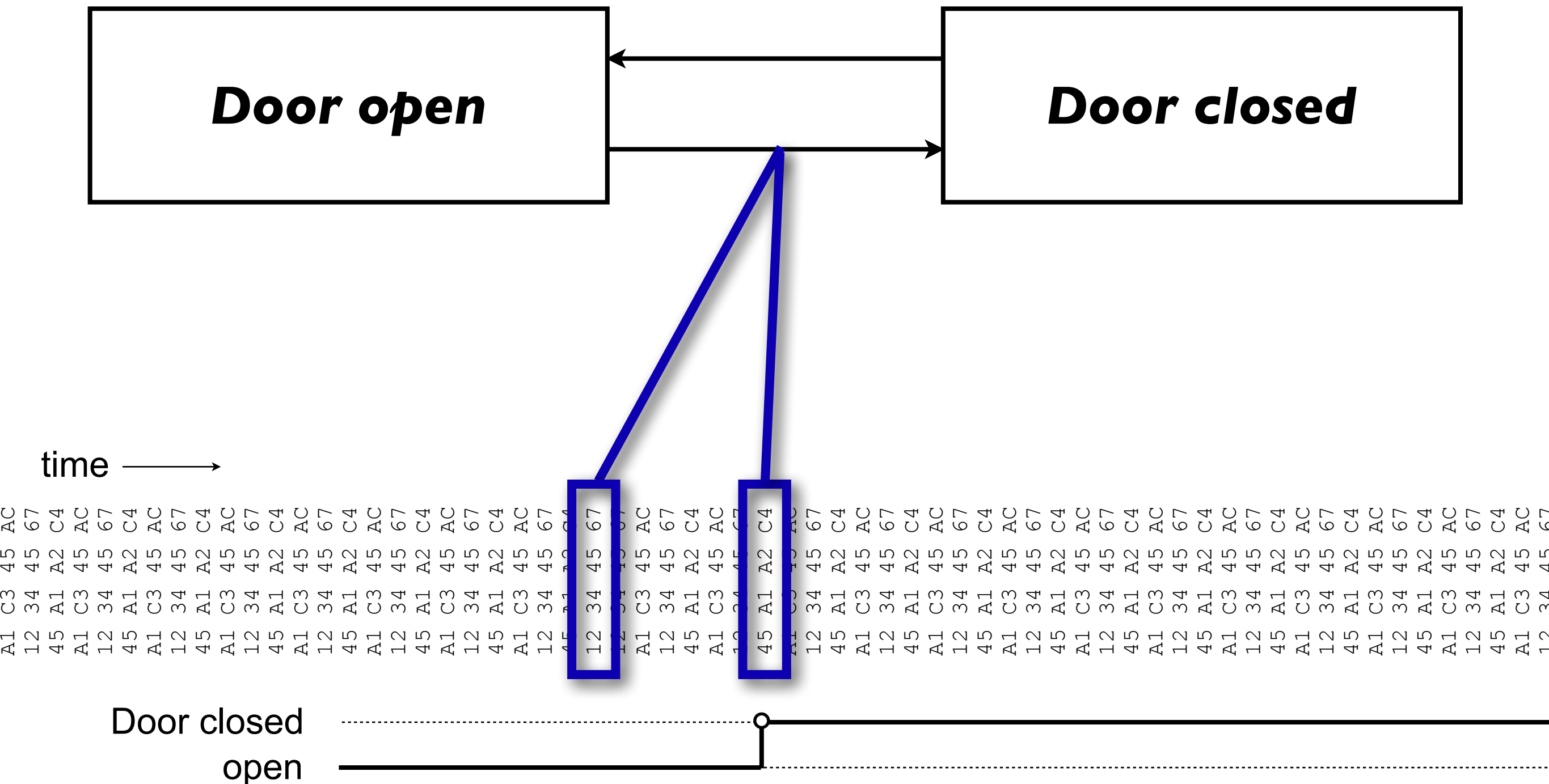


Cardiogram

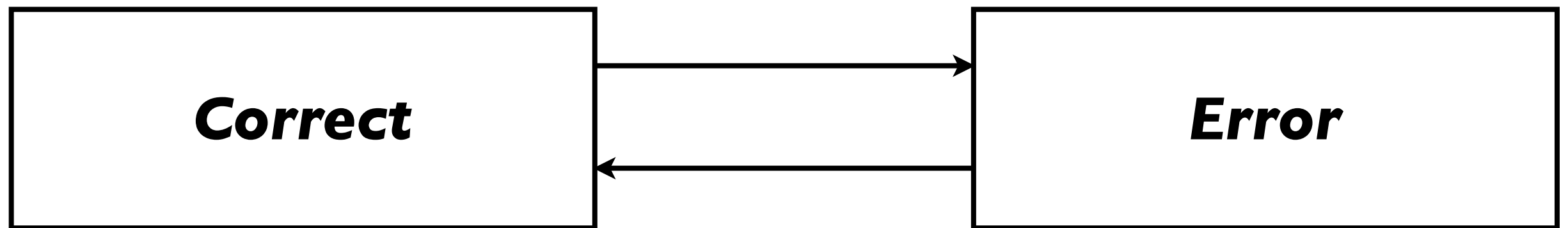
state machine pre-analysis



State Machines: Behavior

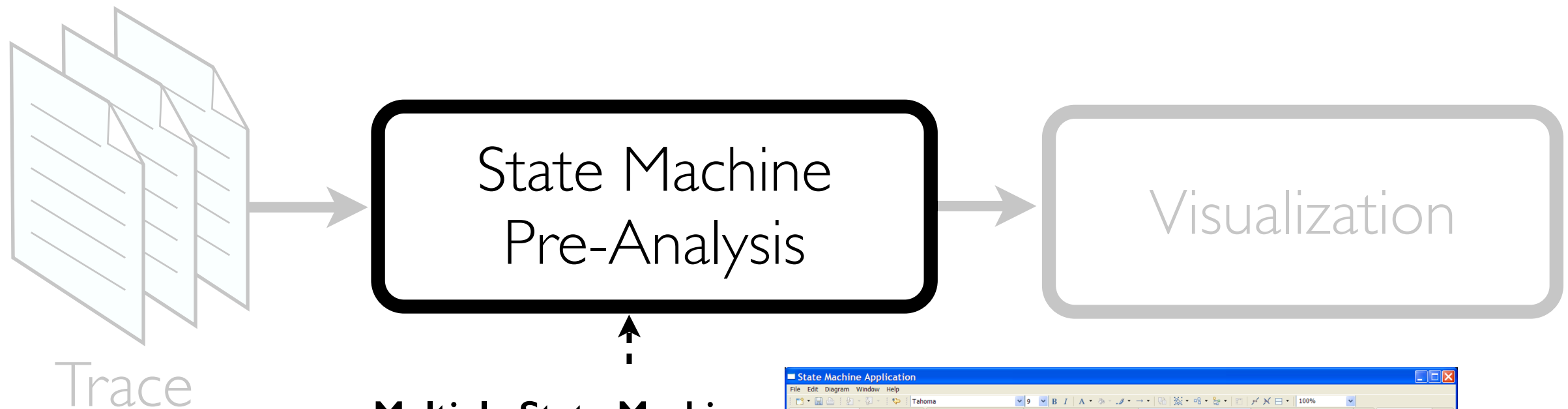


State Machines: Error Detection



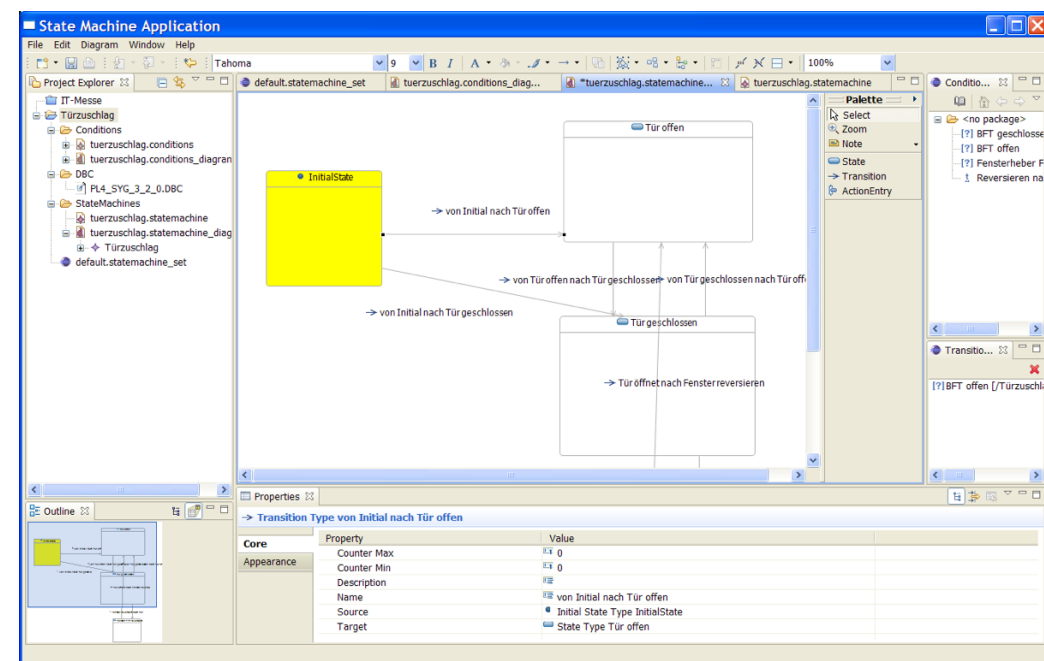
Cardiogram

create/select state machines



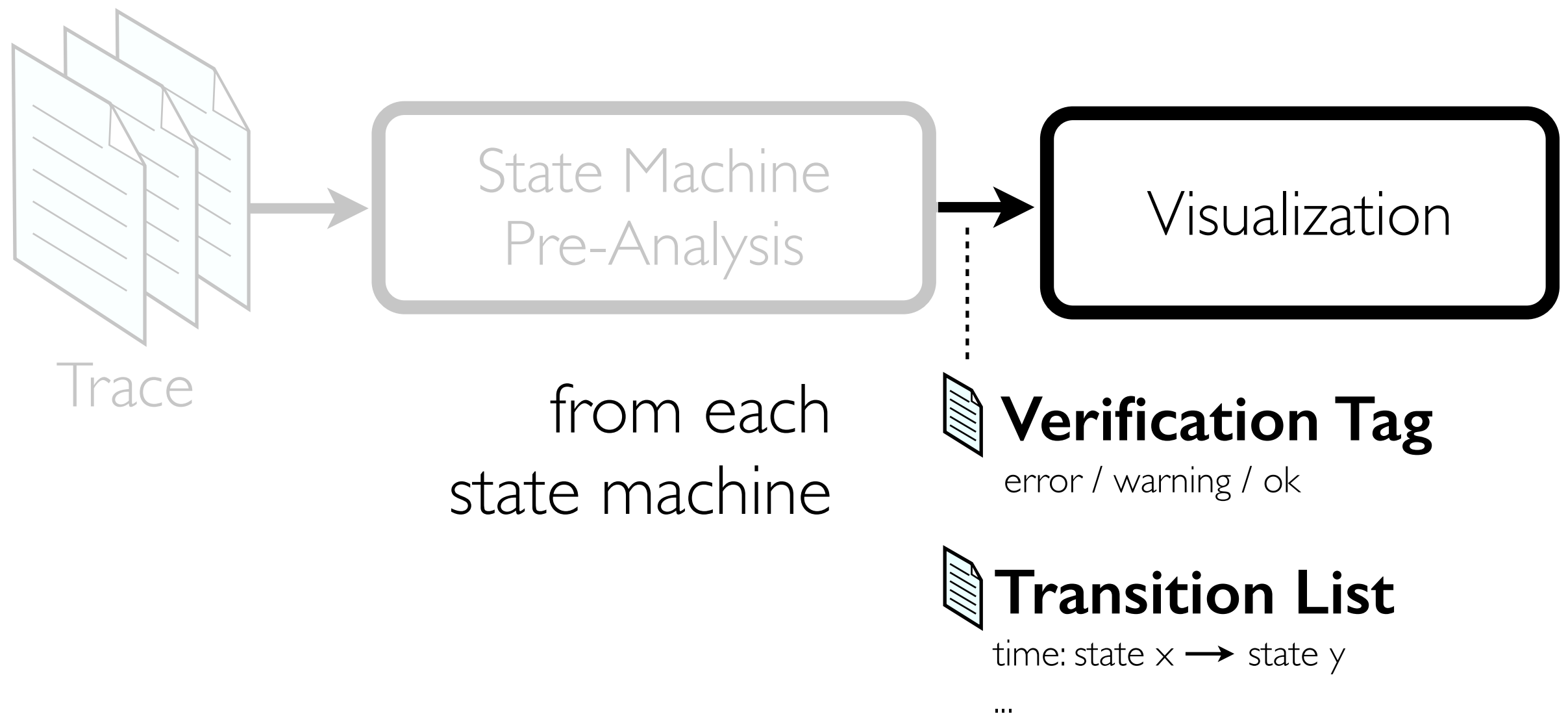
Multiple State Machines

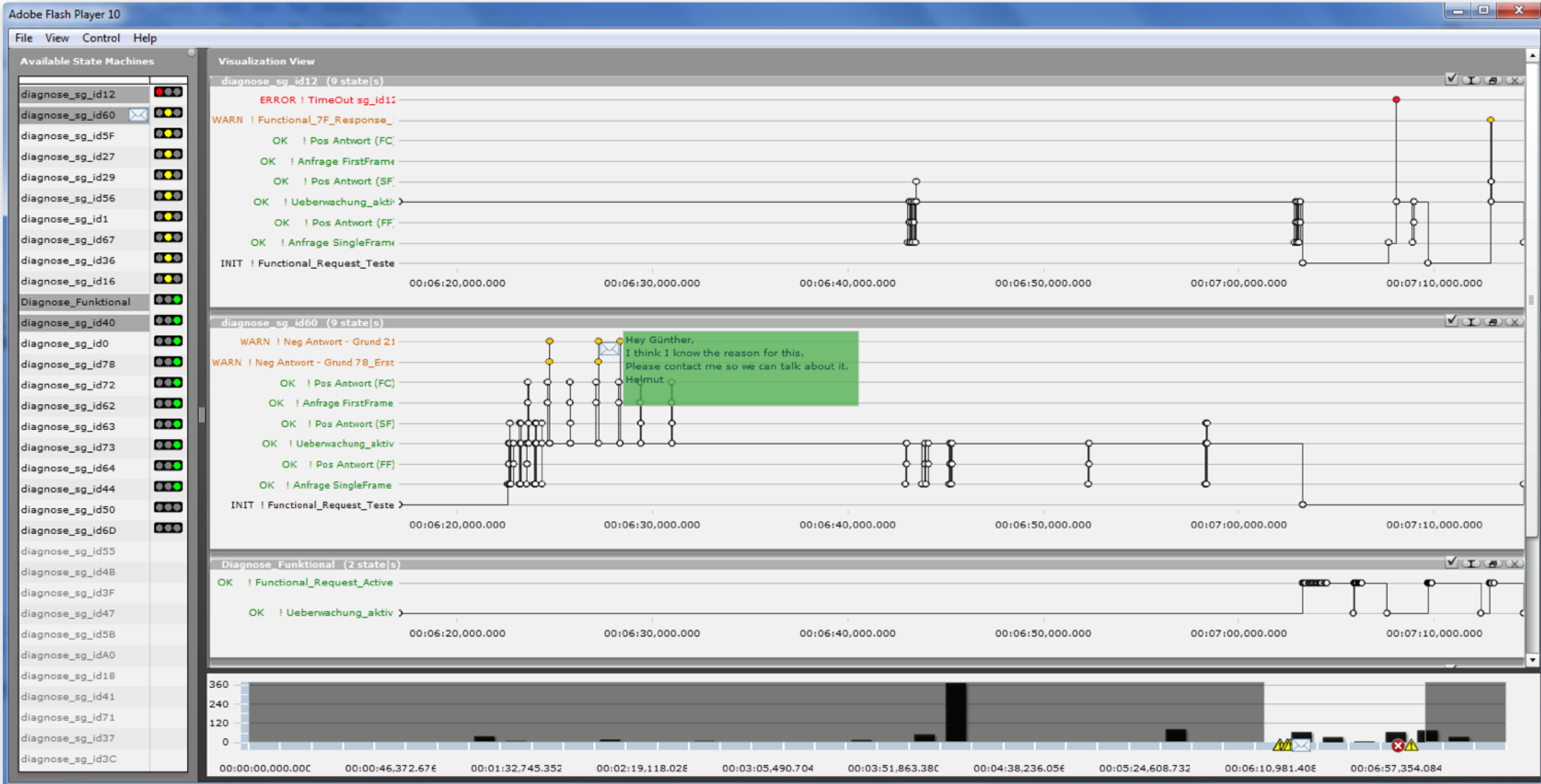
- State Machine 1
- State Machine 2
- State Machine 3
- ... (dozens)

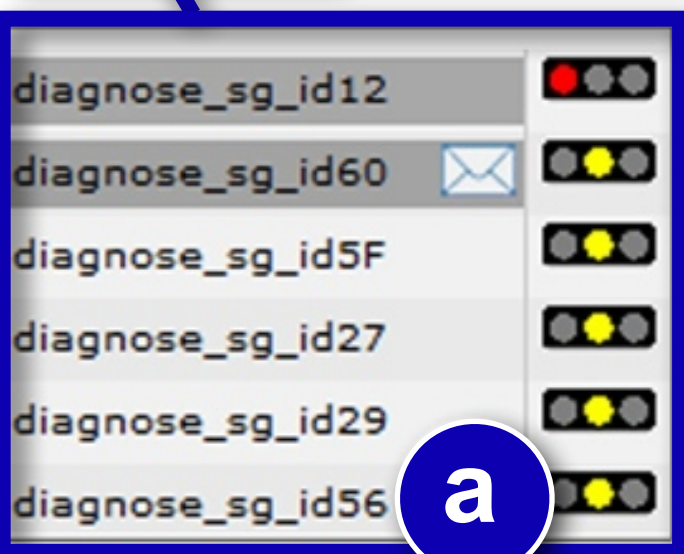
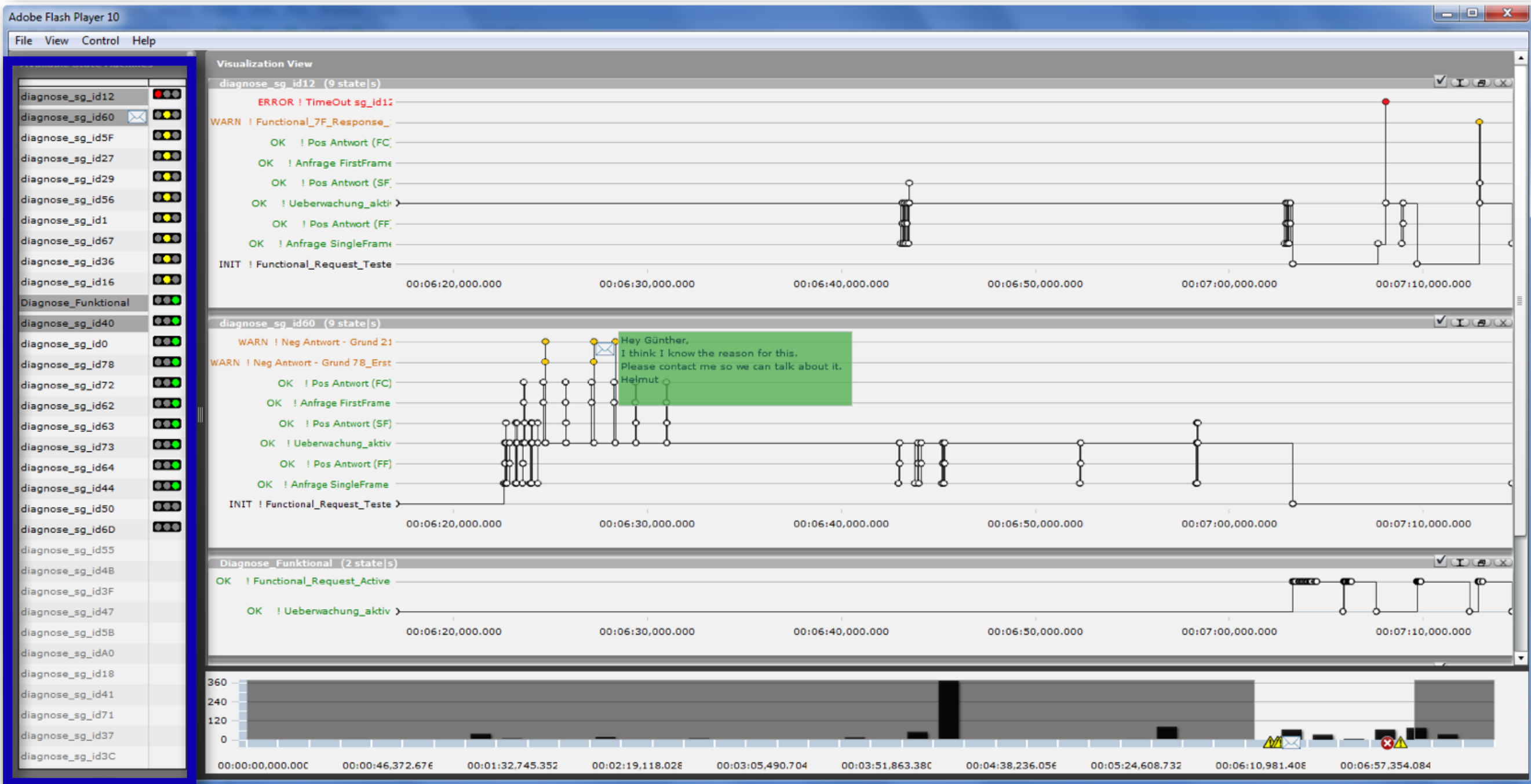


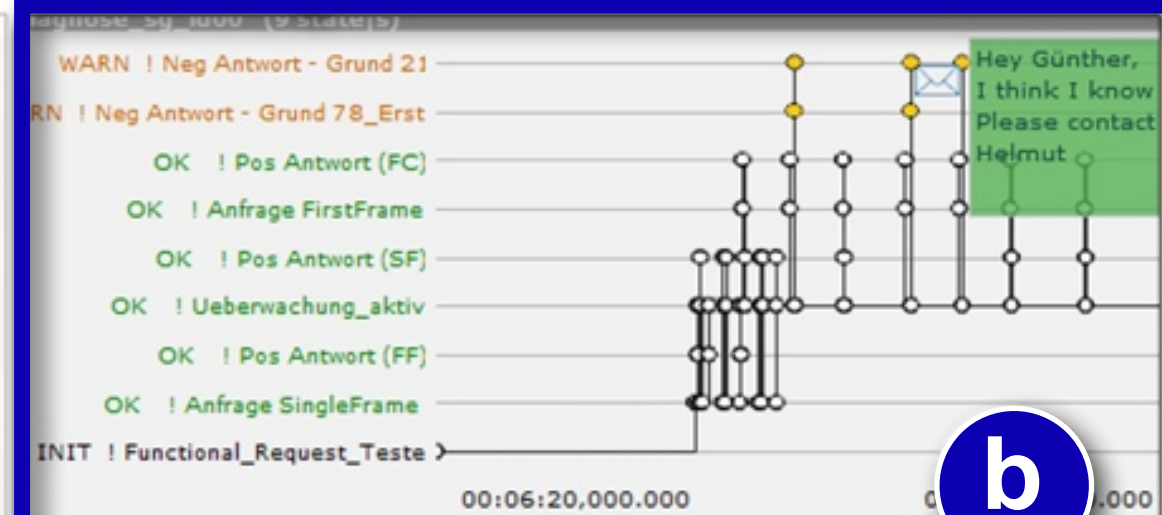
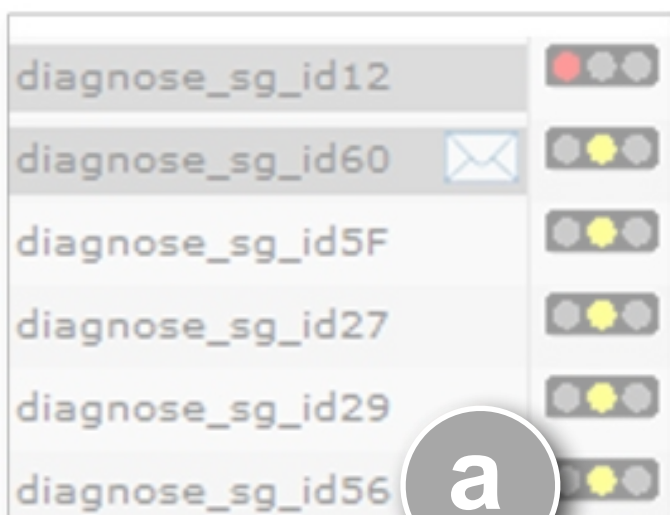
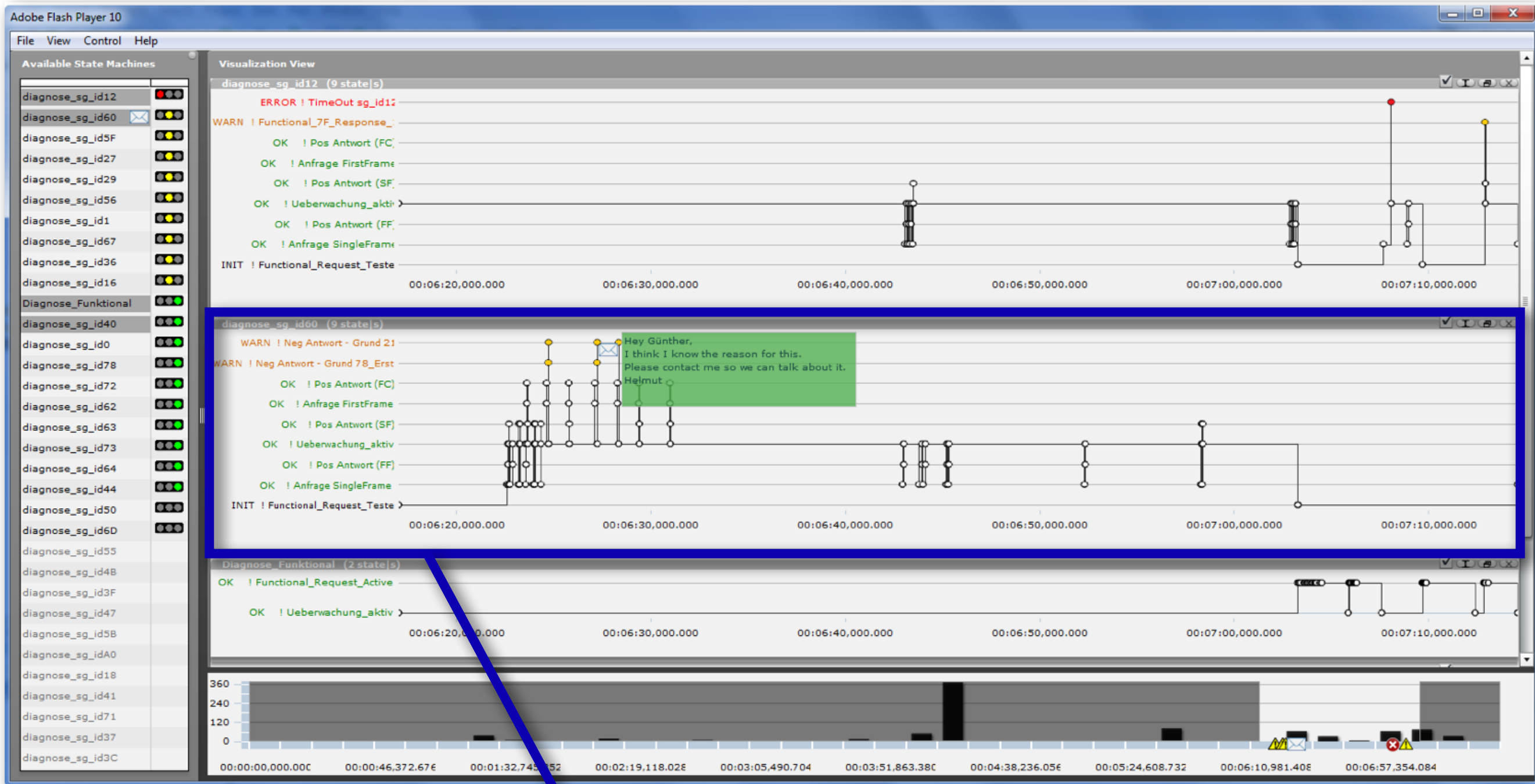
Cardiogram

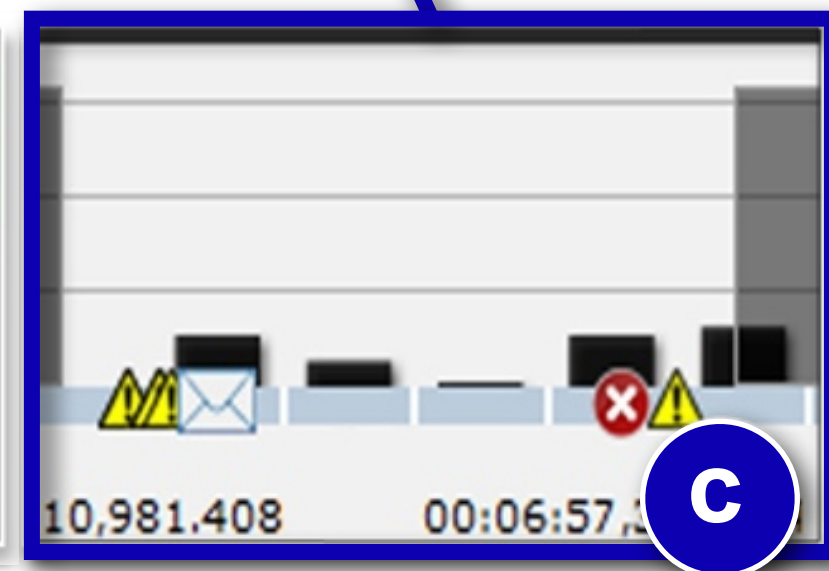
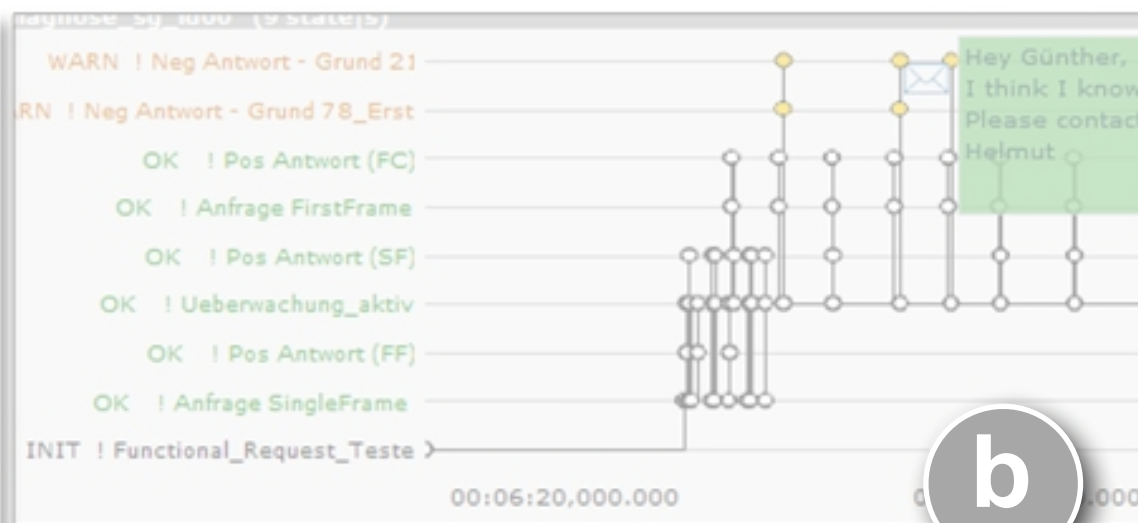
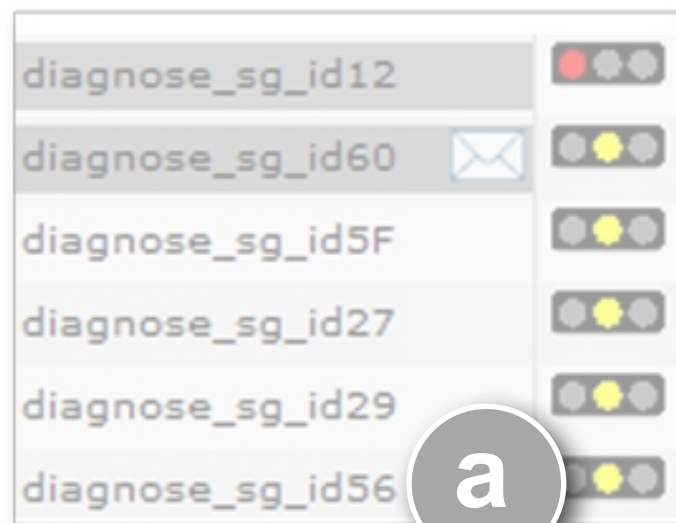
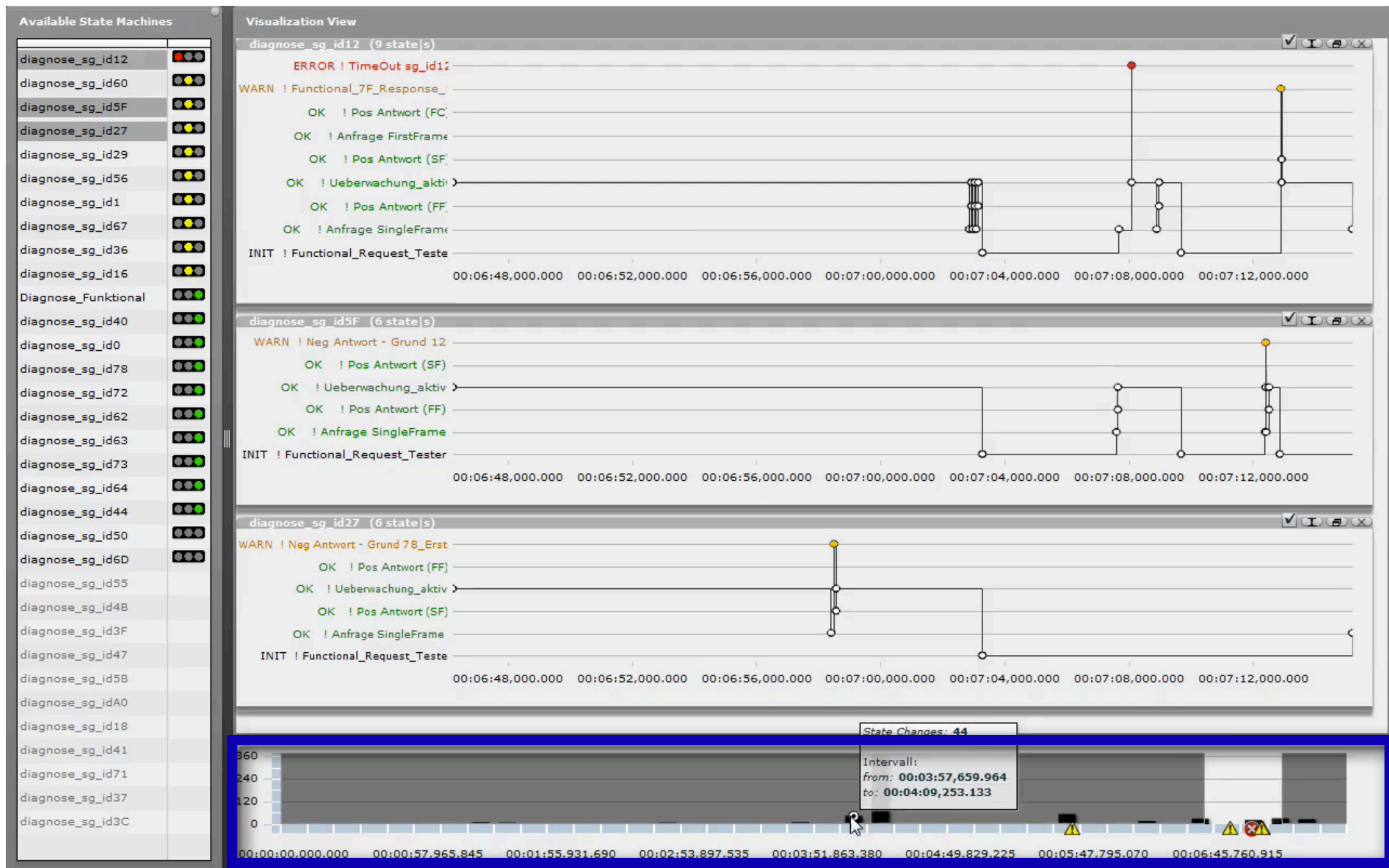
data abstraction/reduction











Longitudinal field study

1 year, 15 engineers

Longitudinal field study results

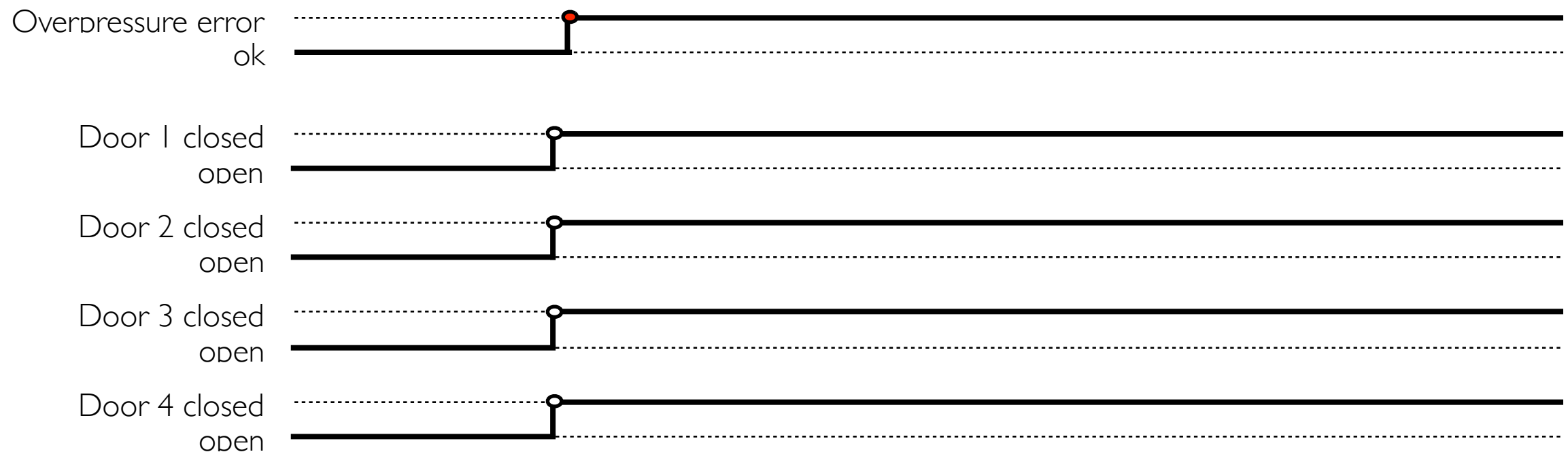
- externalization and sharing of expert knowledge

Longitudinal field study results

- externalization and sharing of expert knowledge
- complete coverage of traces vs. sparse sampling

Longitudinal field study results

- externalization and sharing of expert knowledge
- complete coverage of traces vs. sparse sampling
- understand behavioral correlations



Longitudinal field study results

- externalization and sharing of expert knowledge
- complete coverage of traces vs. sparse sampling
- understand behavioral correlations

deployed and adopted

how to?

“A **design study** is a project in which visualization researchers analyze a specific real-world problem faced by domain experts, design a visualization system that supports solving this problem, validate the design, and reflect about lessons learned in order to refine visualization design guidelines.”

[Sedlmair et al.: Design Study Methodology: Reflections from the Trenches and the Stacks, InfoVis 2012]

design studies:

long and winding road with many pitfalls!



right methods:

Design Study Methodology: Reflection from the Trenches and the Stacks

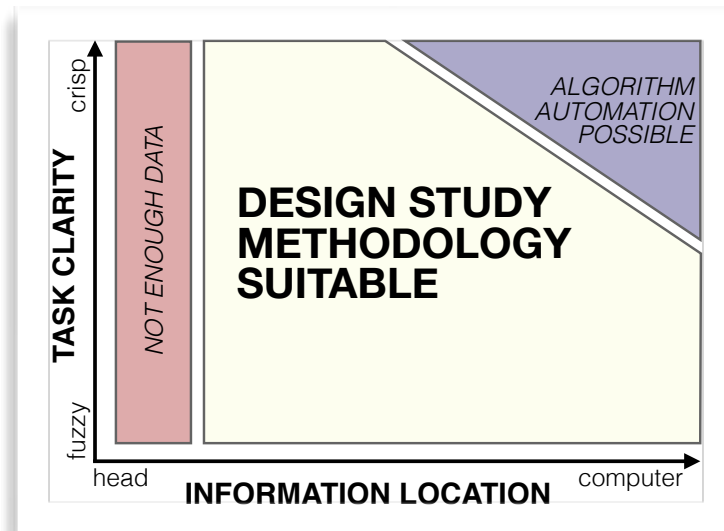
M. Sedlmair, M. Meyer, T. Munzner.

TVCG (Proc. InfoVis 2012). **Best paper honorable mention.**



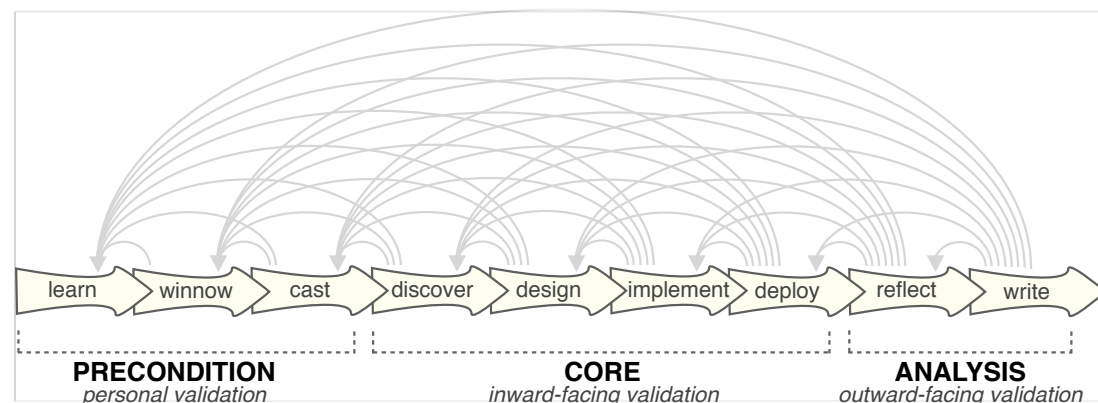
main contributions

guiding other researchers with ...



Definitions

9-stage
framework



PF-1	premature advance: jumping forward over stages	general
PF-2	premature start: insufficient knowledge of vis literature	learn
PF-3	premature commitment: collaboration with wrong people	winnow
PF-4	no real data available (yet)	winnow
PF-5	insufficient time available from potential collaborators	winnow
PF-6	no need for visualization: problem can be automated	winnow
PF-7	researcher expertise does not match domain problem	winnow
PF-8	no need for research: engineering vs. research project	winnow
PF-9	no need for change: existing tools are good enough	winnow

32 pitfalls

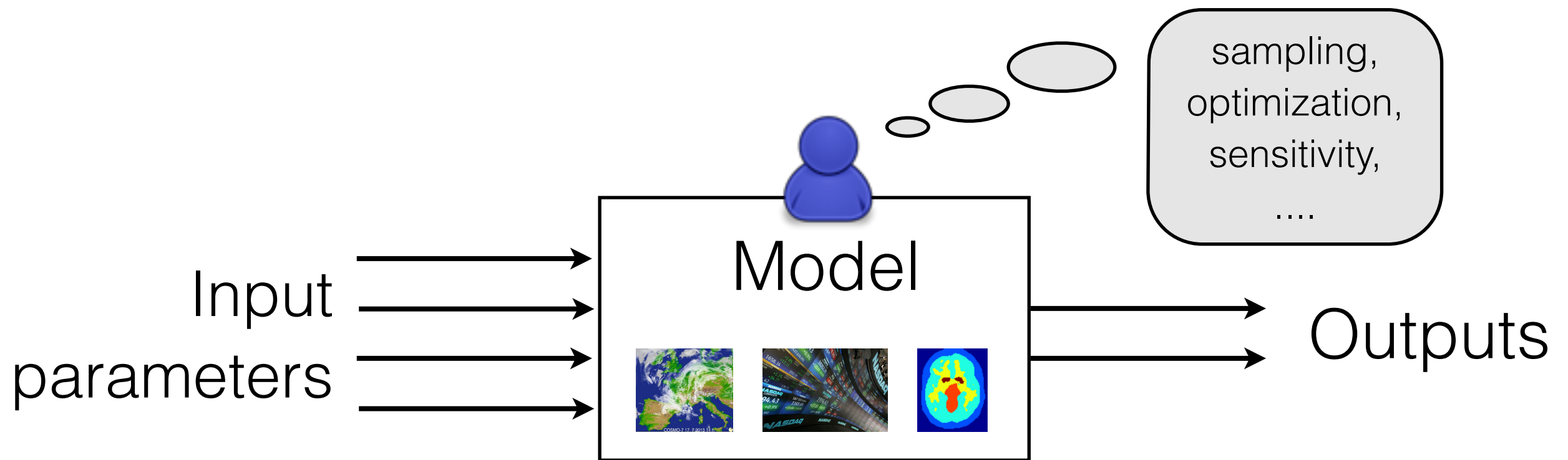
future work

short-term

1. Design studies

- real users/real data:
 - contribute to their fields
 - improve our guidelines
- right methods:
 - validate, improve, extend design study methodology

Design studies: **Visual Parameter Space Analysis**



Sedlmair, Heinzl, Bruckner, Piringner, Möller: *Visual Parameter Space Analysis: A Conceptual Framework*, InfoVis 2014

Bergner, **Sedlmair**, Möller, Abdolyousef, Saad: *ParaGlide: Interactive Parameter Space Partitioning for Computer Simulations*, IEEE TVCG, 2013

Design studies

Data-driven Journalism



- *users:* journalists
- *tasks:* story creation
- *data:* text over time, e.g., parliament speeches



FFG-funded VALiD project (690.000 Euro)

Aigner (PI), **Sedlmair (Co-PI)**, Ausserhofer (Co-PI), Popper (Co-PI)

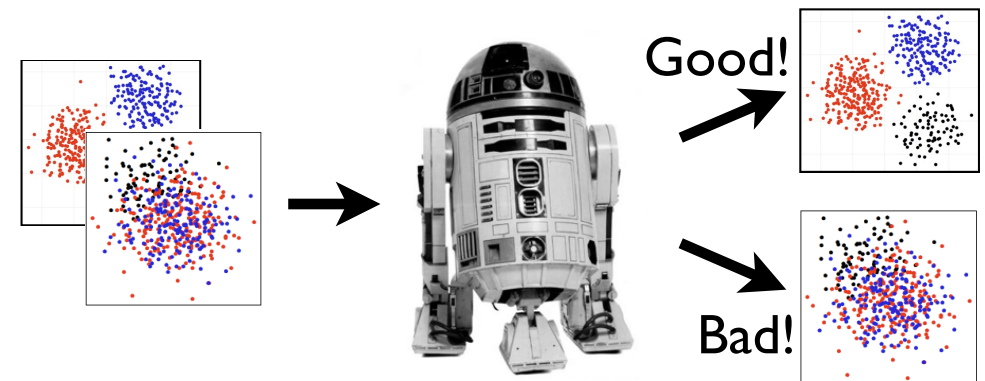
2. High-dim data analysis techniques

- real users/real data:
 - understand techniques in real circumstances
 - build better techniques
- right method
 - combine approaches from HCI, Vis, Psychology, and Machine Learning

High-dim techniques: **visual quality measures**

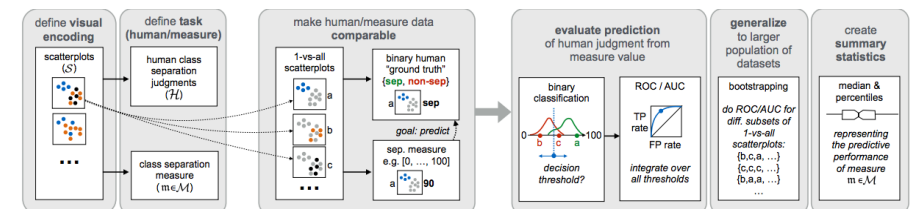
- real users/real data

[**Sedlmair**, Tatu, Munzner, Tory: *A Taxonomy of Visual Cluster Separation Factors*, EuroVis 2012]



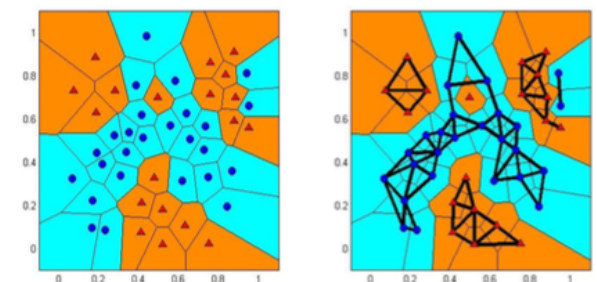
- right method

[**Sedlmair**, Aupetit, *Data-driven Evaluation of Visual Quality Measures*, EuroVis 2015]



- new techniques

[Aupetit, **Sedlmair**: *Learning human visual judgment of class separation in colored scatterplots for Visual Analytics*, NIPS 2015 - in submission]



High-dim techniques: **other ...**

- clustering
- dimension reduction
- sensitivity analysis
- ...



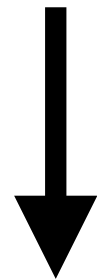
future work

long-term

Vision:

The “cookbook” for Visual Data Analysis

abstract data & task



data analysis technique



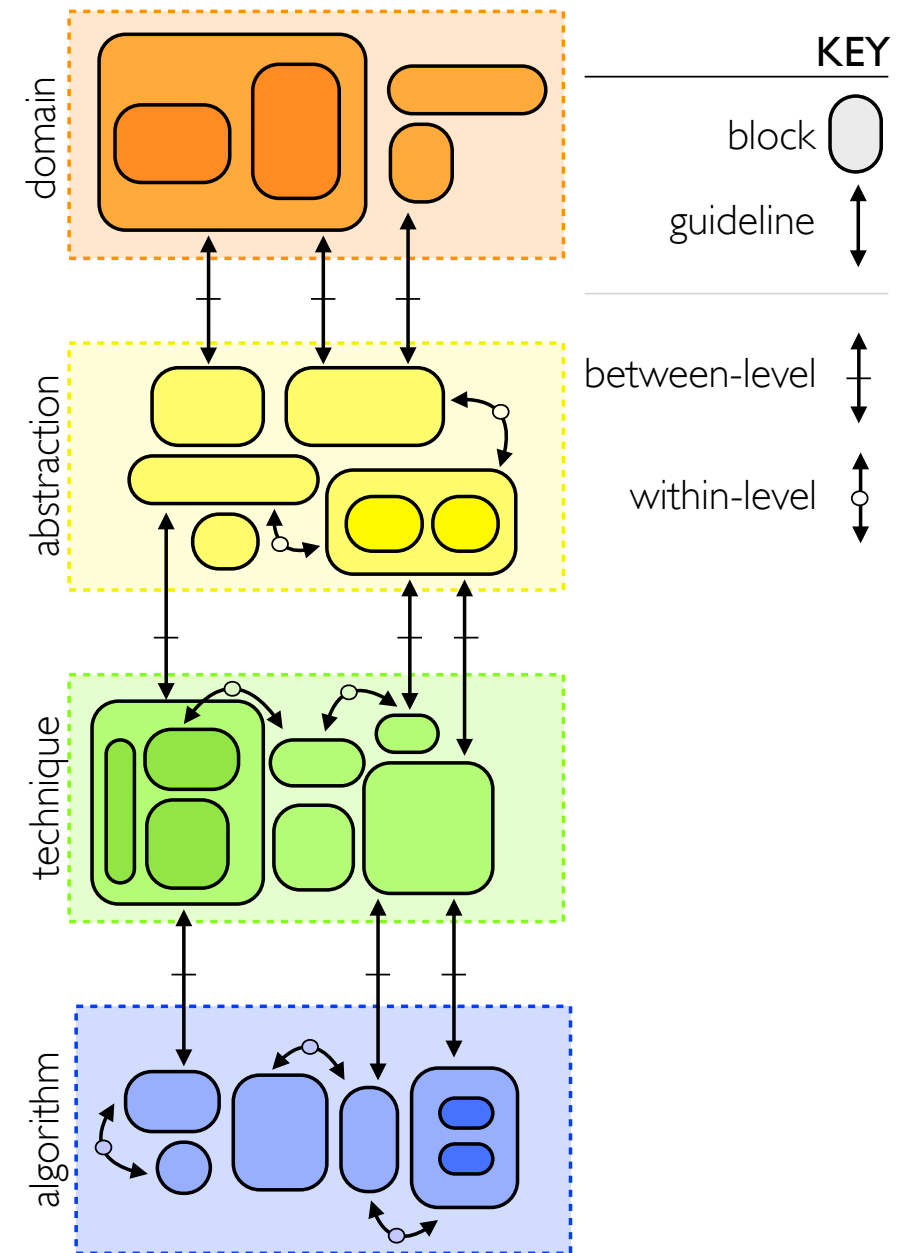
The “cookbook” for Visual Data Analysis

blocks and guidelines

Meyer, **Sedlmair**, Quinan, Munzner:
The Nested Blocks and Guidelines Model,
IVS Journal 2013.

design decision patterns

ongoing project:
with Heidi Lam & Miriah Meyer



The “cookbook” for Visual Data Analysis

natural science
vs. design research

vast amounts of
real users/real data
work necessary



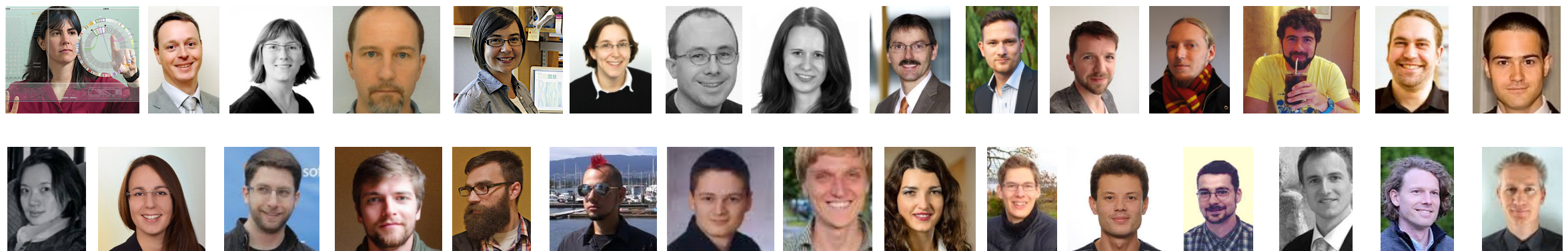
[Sir Francis Bacon 1561-1626]

Thanks to you & to ...

VDA Group @ University of Vienna



Many, many other collaborators



...

Questions?

On CS Undergrad Teaching and Visualization Research

1. Sorting in different programming paradigms
2. Visual Data Analysis:
real users, real data, right methods

Michael Sedlmair

Visualization and Data Analysis group (VDA)



universität
wien